

NVRAMOS '14

10.30. 2014

Resolving Journaling of Journal Anomaly via Weaving Recovery Information into DB Page

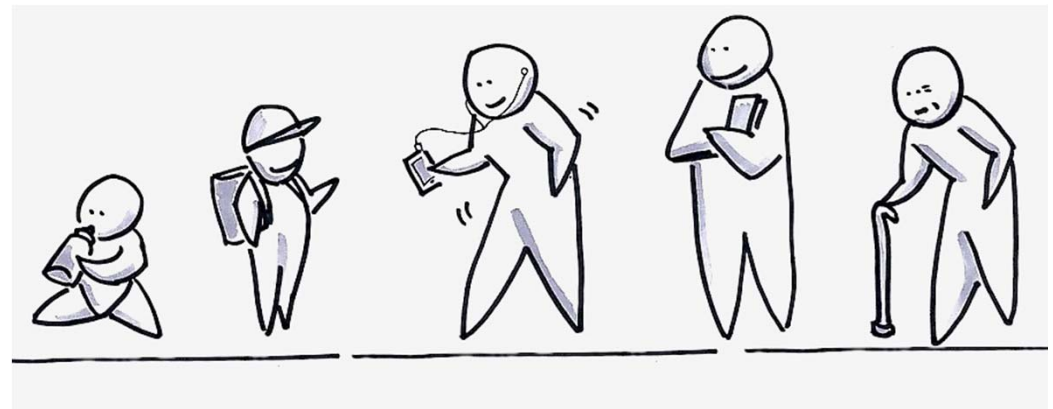
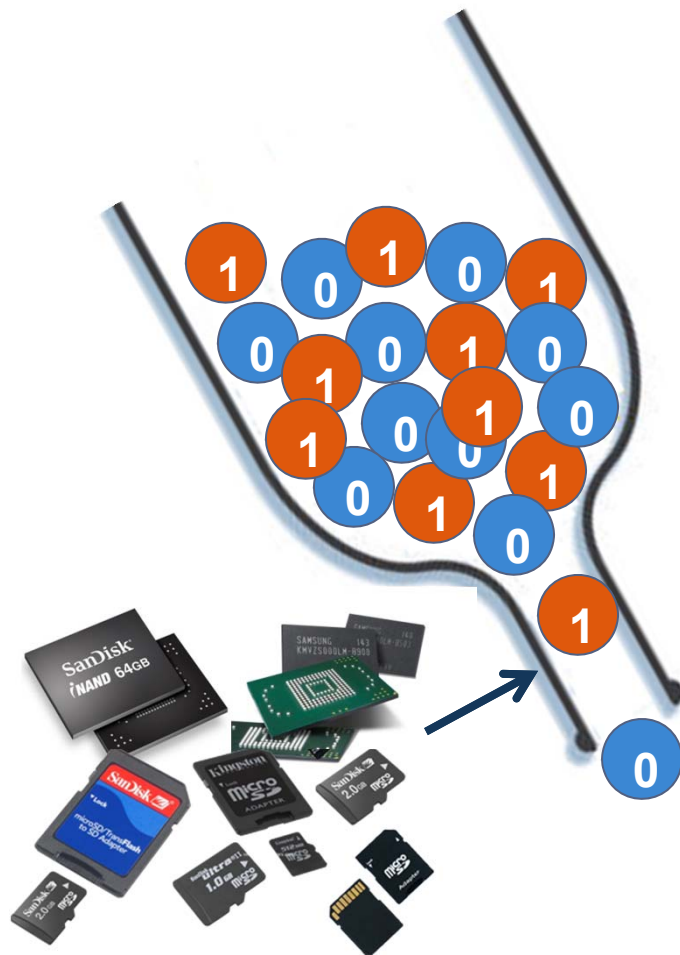
Beomseok Nam
UNIST

Outline

- Motivation
 - Journaling of Journal Anomaly
- How to resolve Journaling of Journal anomaly
 - Multi-Version B-Tree (MVBT)
- Optimizations of Multi-Version B-Tree
 - Lazy Split
 - Reserved Buffer Space
 - Lazy Garbage Collection
 - Metadata Embedding
 - Disabling Sibling Redistribution
- Evaluation
- Conclusion

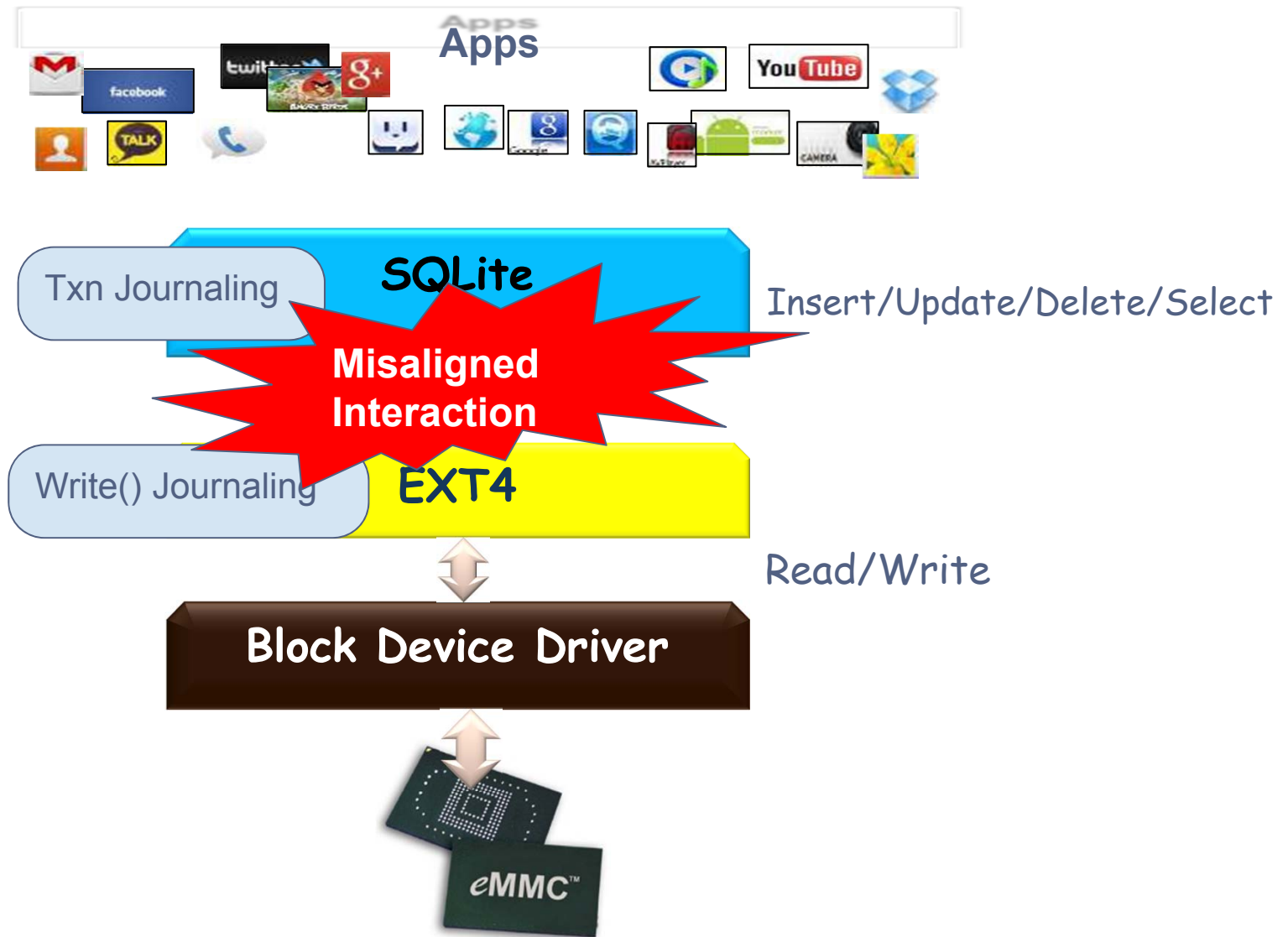
Storage I/O Problems in Android

- Performance Bottleneck
- Lifetime of Storage



Cause = Excessive IO

Android I/O Stack

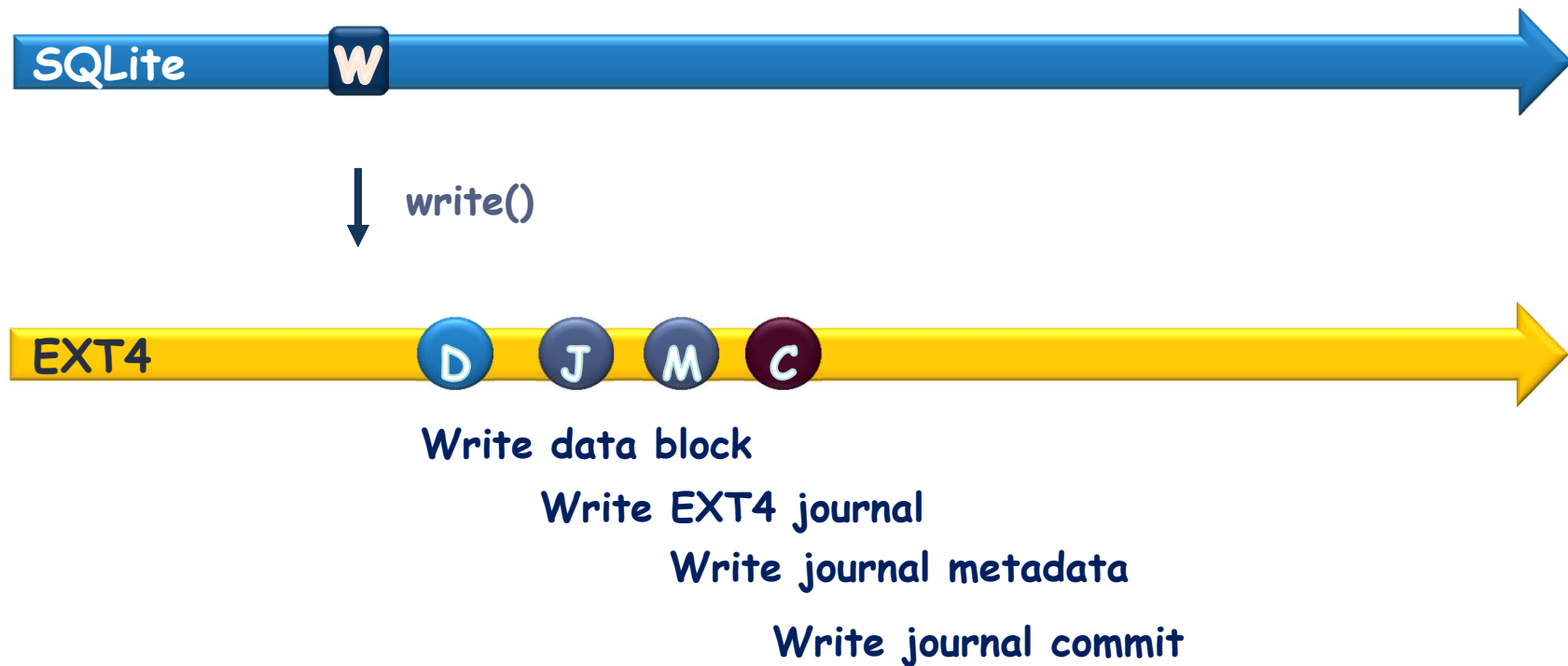


Journaling of Journal Anomaly

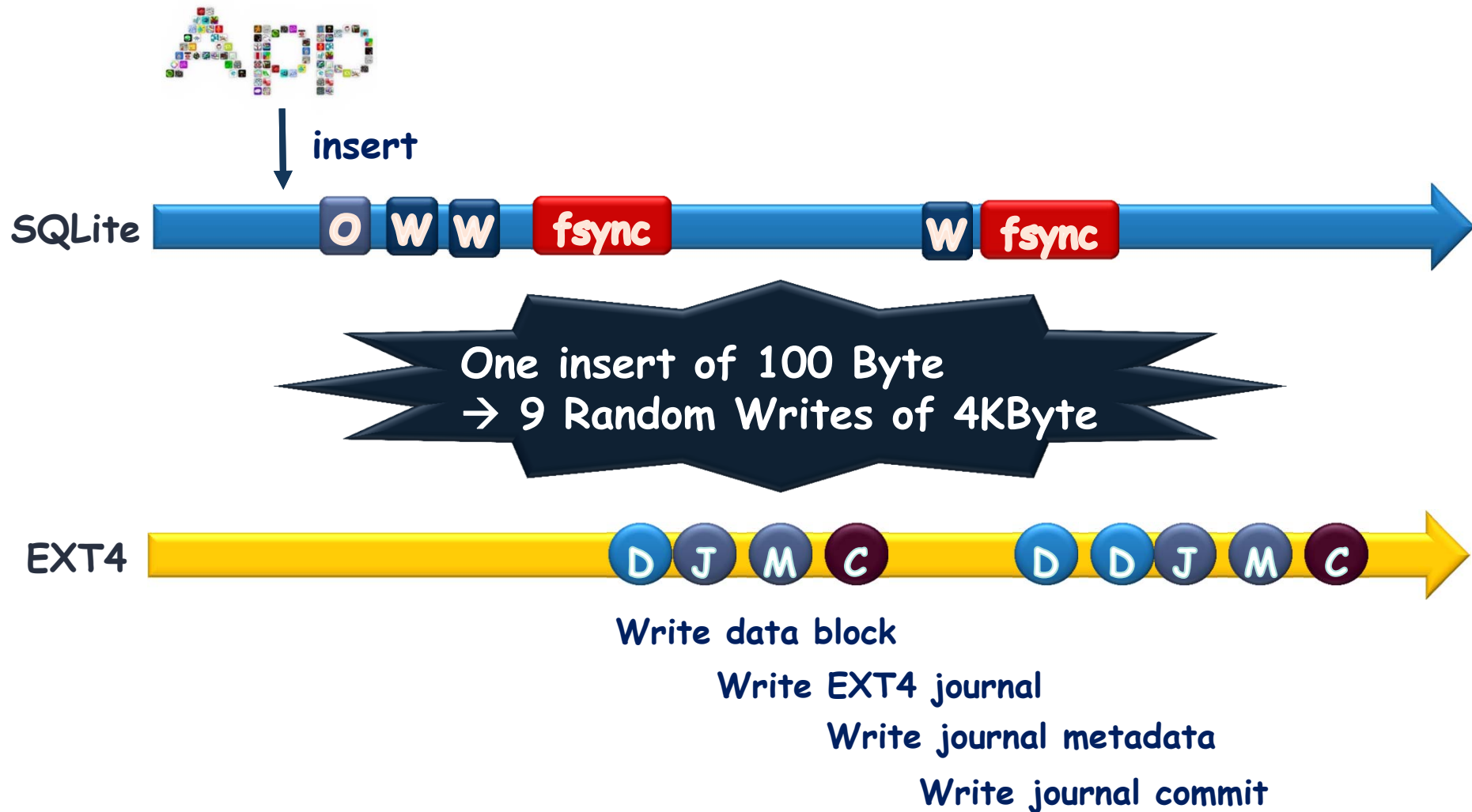
- Journaling in SQLite (TRUNCATE mode)



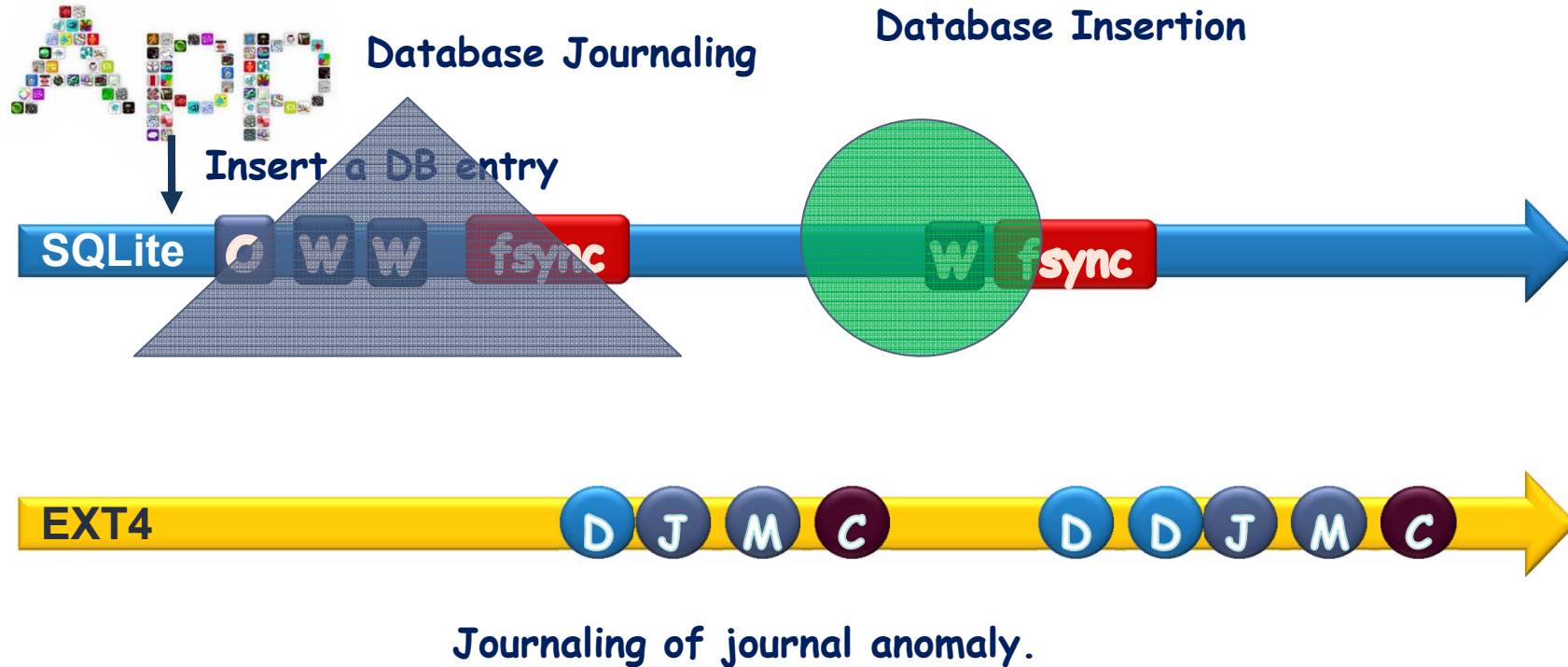
EXT4 Journaling (ordered mode)



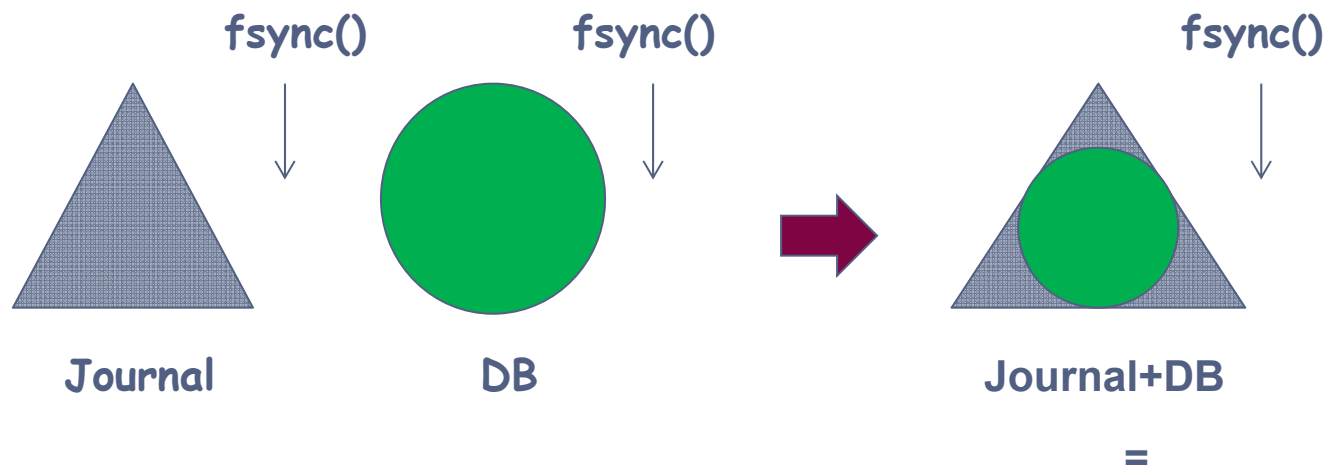
Journaling of Journal Anomaly



How to Resolve Journaling of Journal?



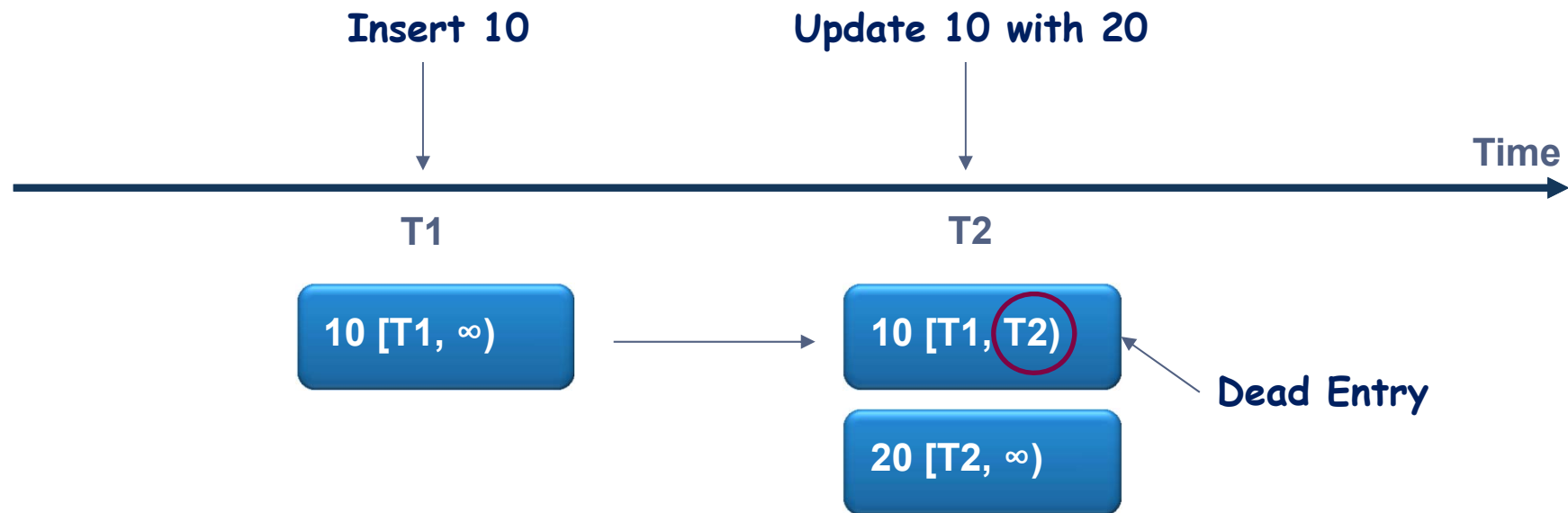
How to Resolve Journaling of Journal?



**Version-based B-Tree
(MVBT)**

Version-based B-Tree (MVBT)

■ Versioning



Do not overwrite old version data
Do not need a rollback journal

Node Split in Multi-Version B-Tree

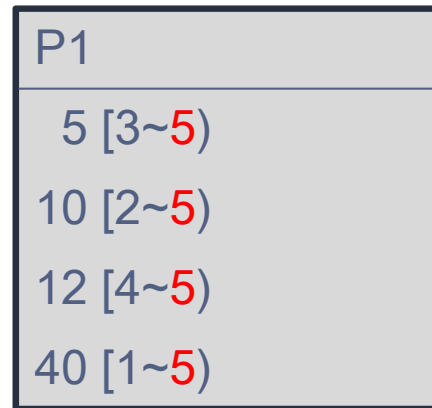
key 25, ver=5



P1
5 [3~ ∞)
10 [2~ ∞)
12 [4~ ∞)
40 [1~ ∞)

Node Split in Multi-Version B-Tree

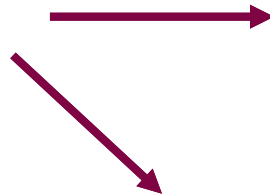
key 25, ver=5



Dead
Node

Node Split in Multi-Version B-Tree

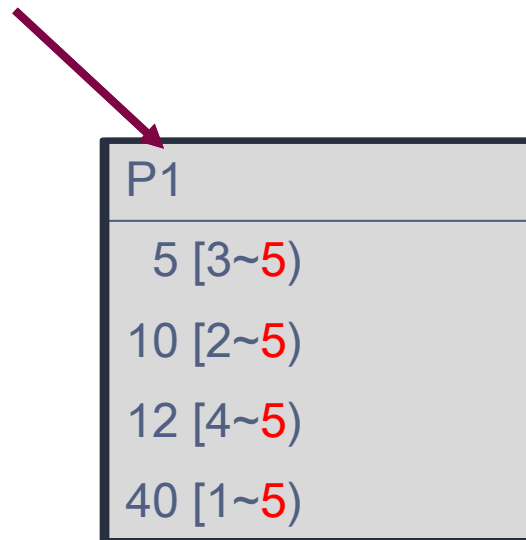
key 25, ver=5



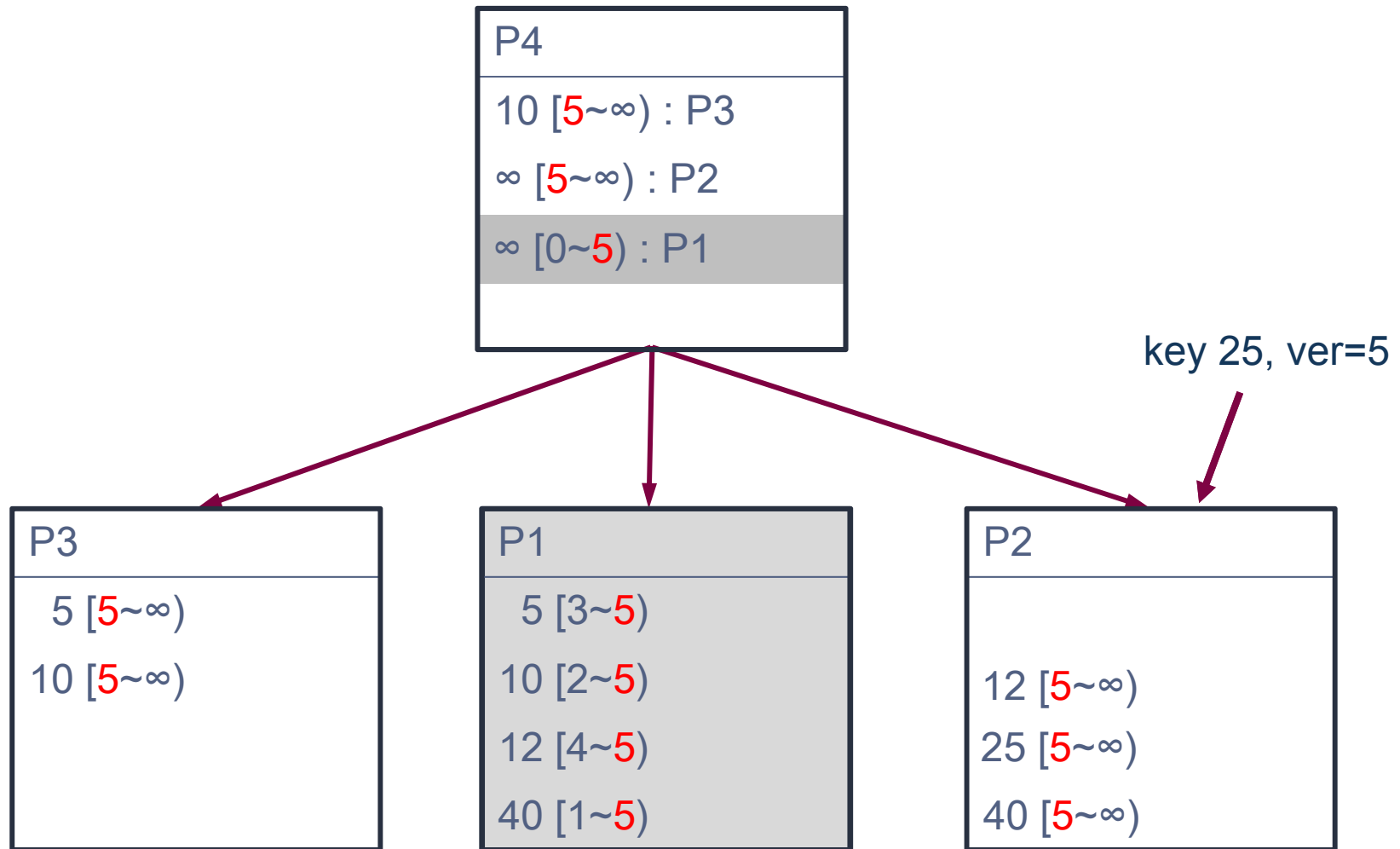
P1
5 [3~5)
10 [2~5)
12 [4~5)
40 [1~5)

Node Split in Multi-Version B-Tree

key 25, ver=5

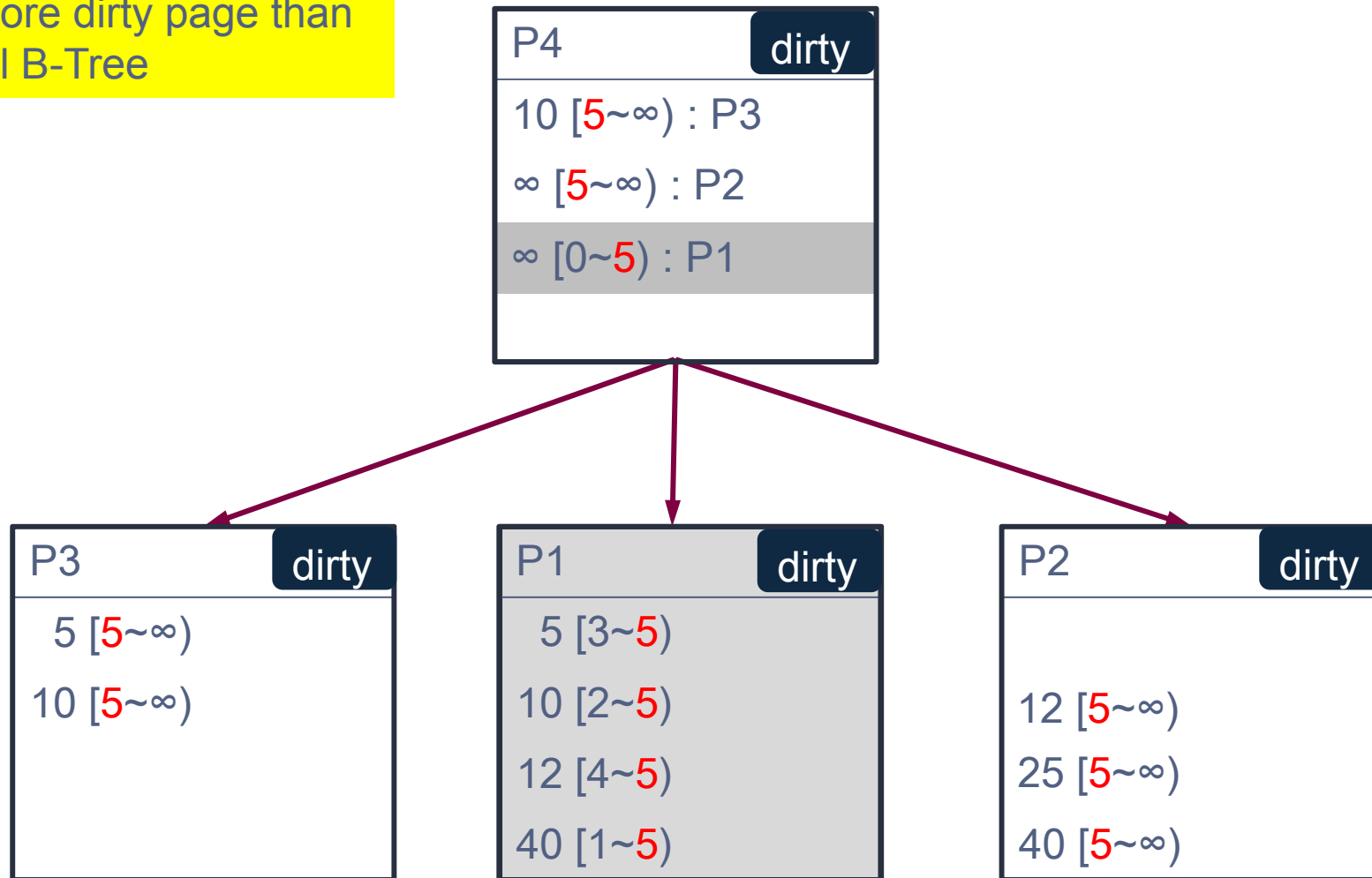


Node Split in Multi-Version B-Tree

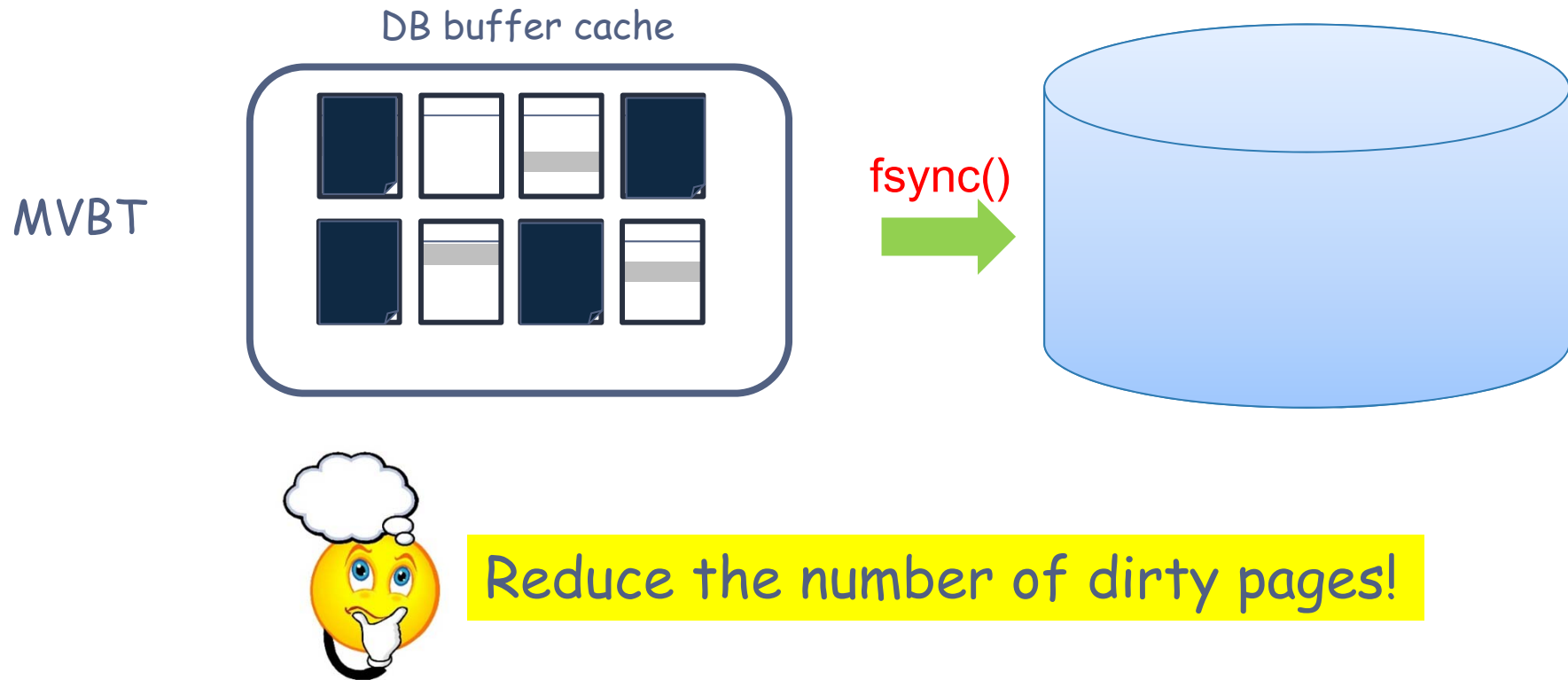


Node Split in Multi-Version B-Tree

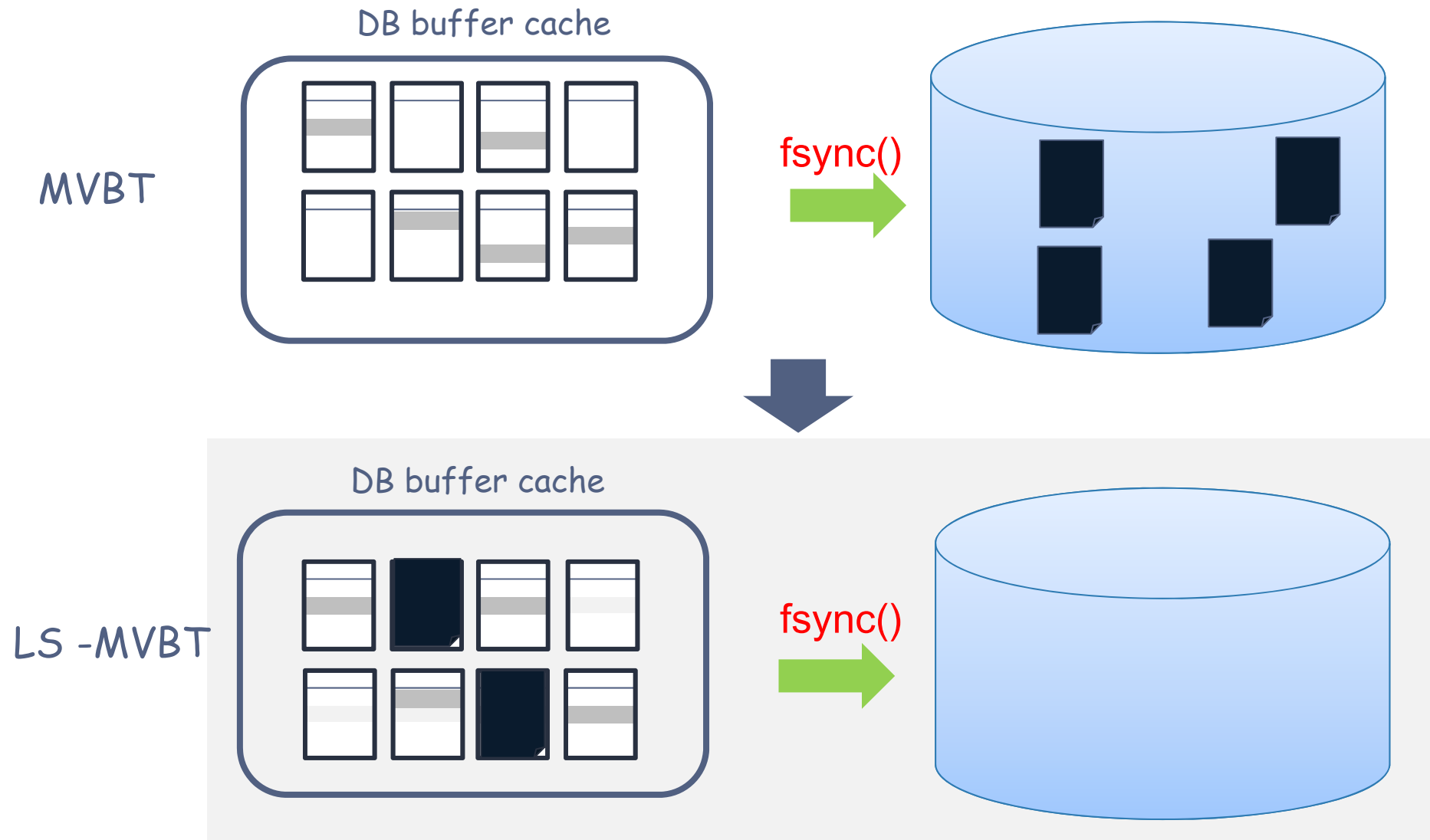
One more dirty page than original B-Tree



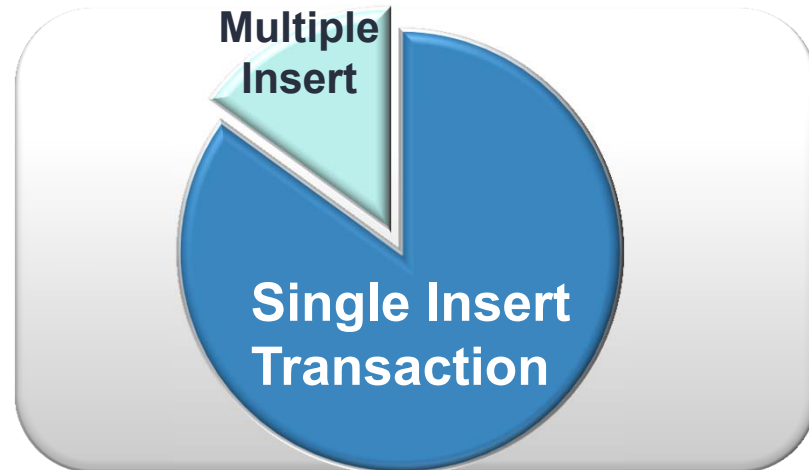
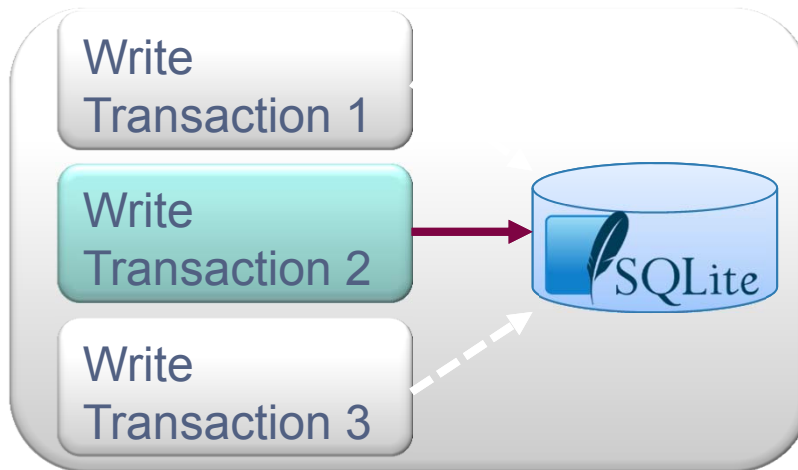
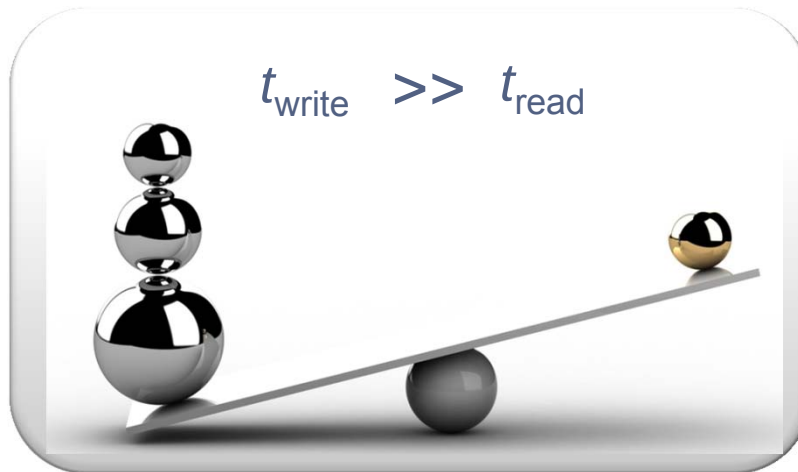
I/O Traffic in MVBT



I/O Traffic in MVBT

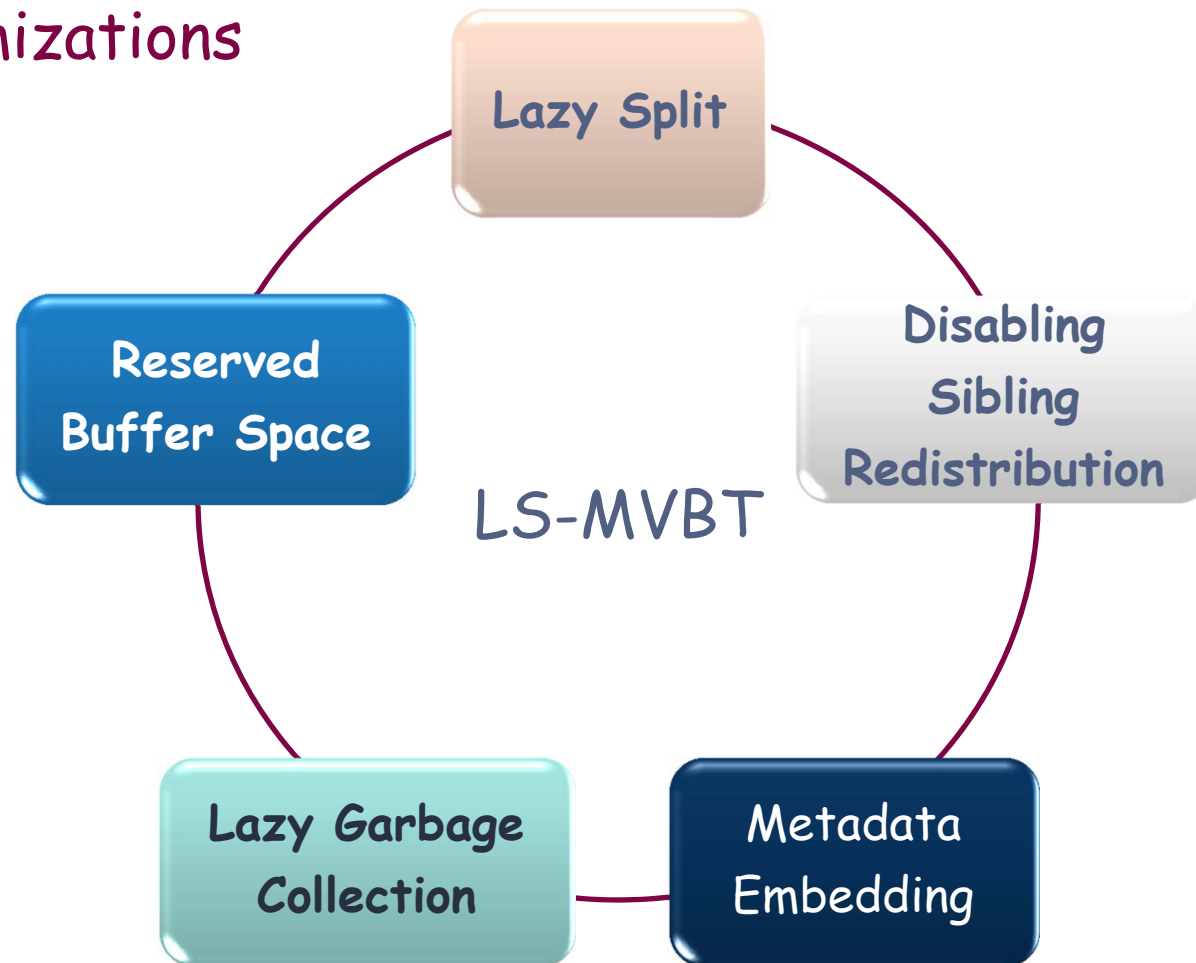


Optimizations in Android I/O



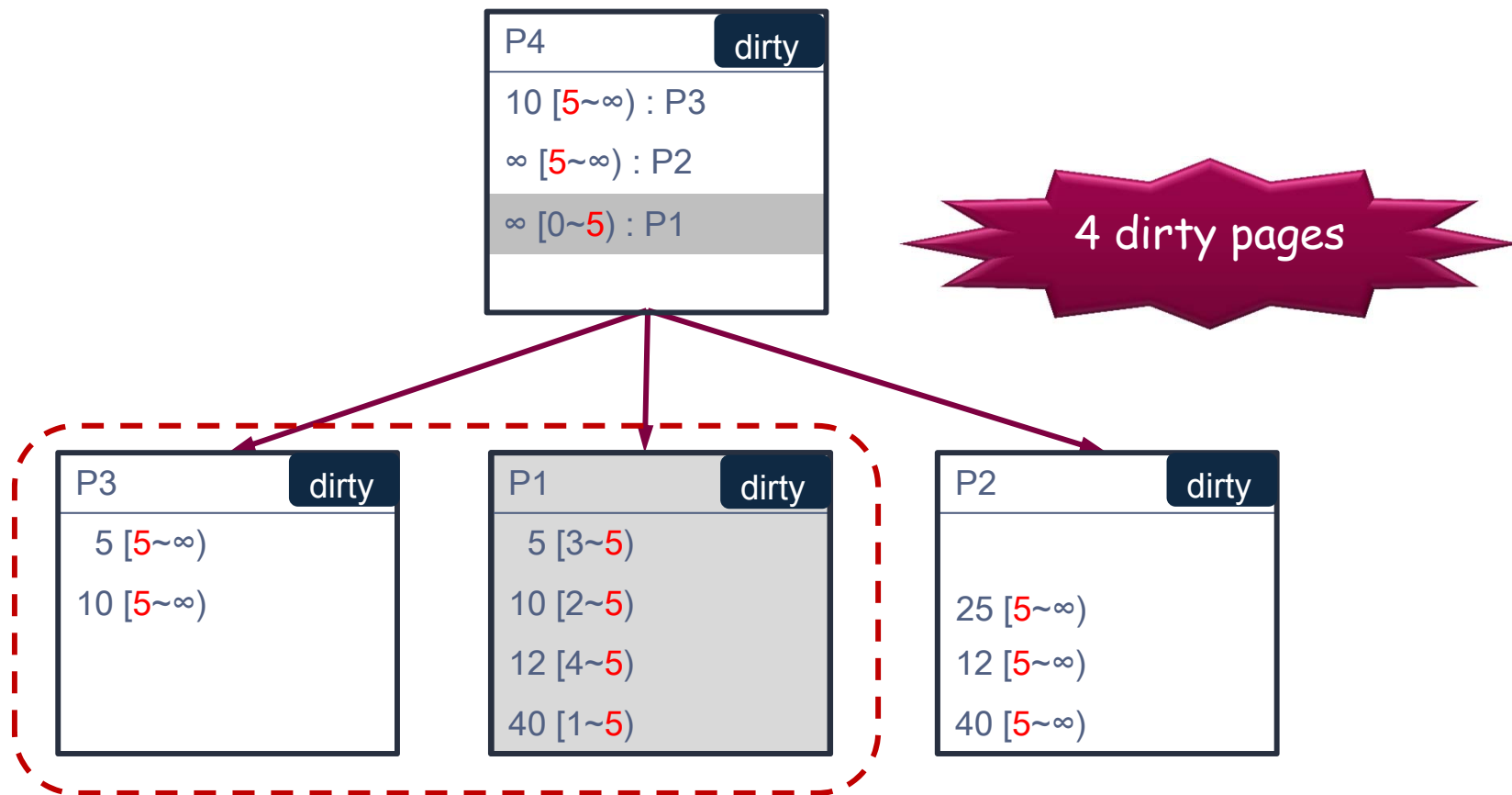
Lazy Split Multi-Version B-Tree (LS-MVBT)

- Optimizations



Optimization1: Lazy Split

- Legacy Split in MVBT



Optimization1: Lazy Split

key 25, ver=5



P1
5 [3~ ∞)
10 [2~ ∞)
12 [4~ ∞)
40 [1~ ∞)

Optimization1: Lazy Split

key 25, ver=5

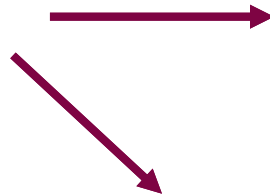


P1
5 [3~ ∞)
10 [2~ ∞)
12 [4~5)
40 [1~5)

Lazy Node:
Half-dead,
Half-live

Optimization1: Lazy Split

key 25, ver=5



P1
5 [3~ ∞)
10 [2~ ∞)
12 [4~5)
40 [1~5)

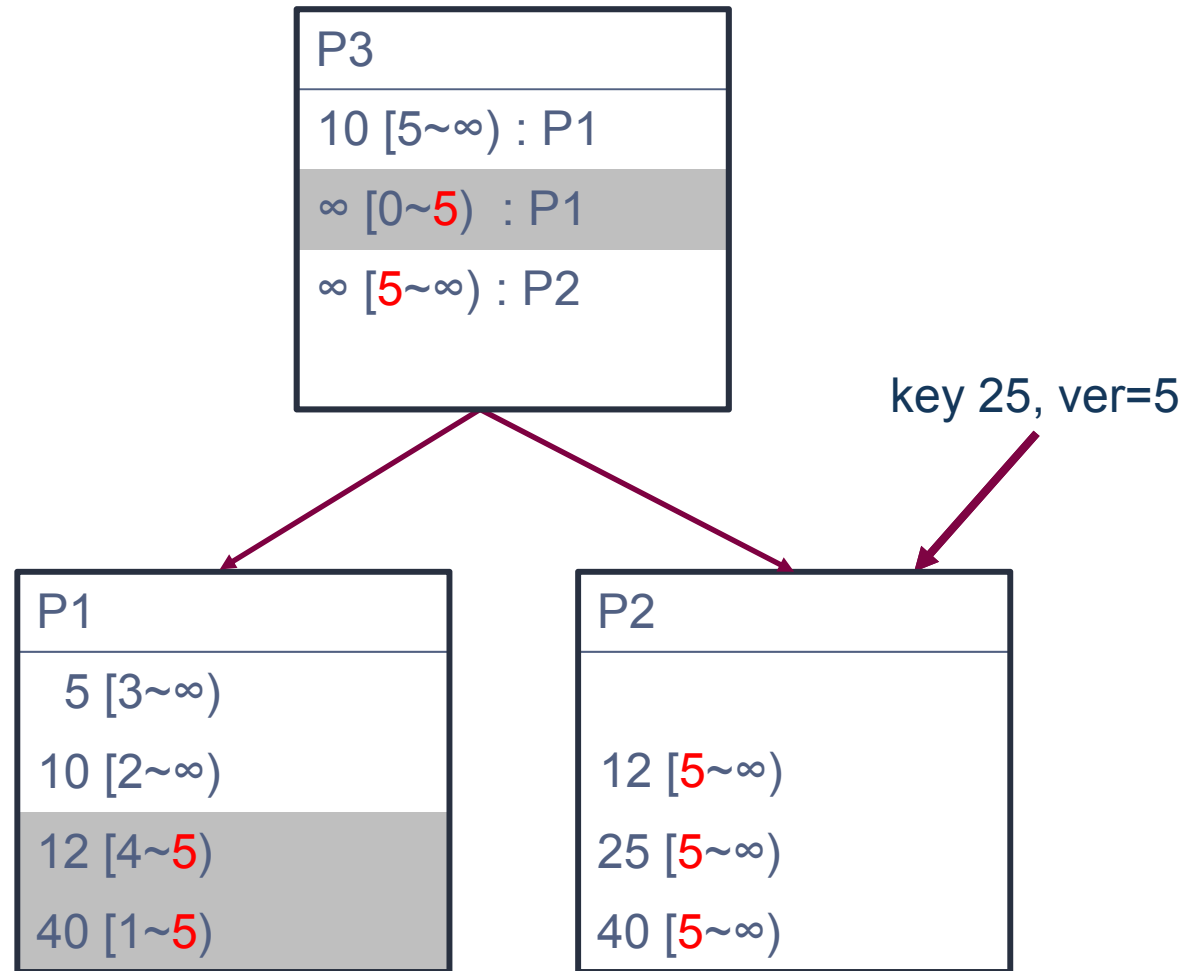
Optimization1: Lazy Split

key 25, ver=5



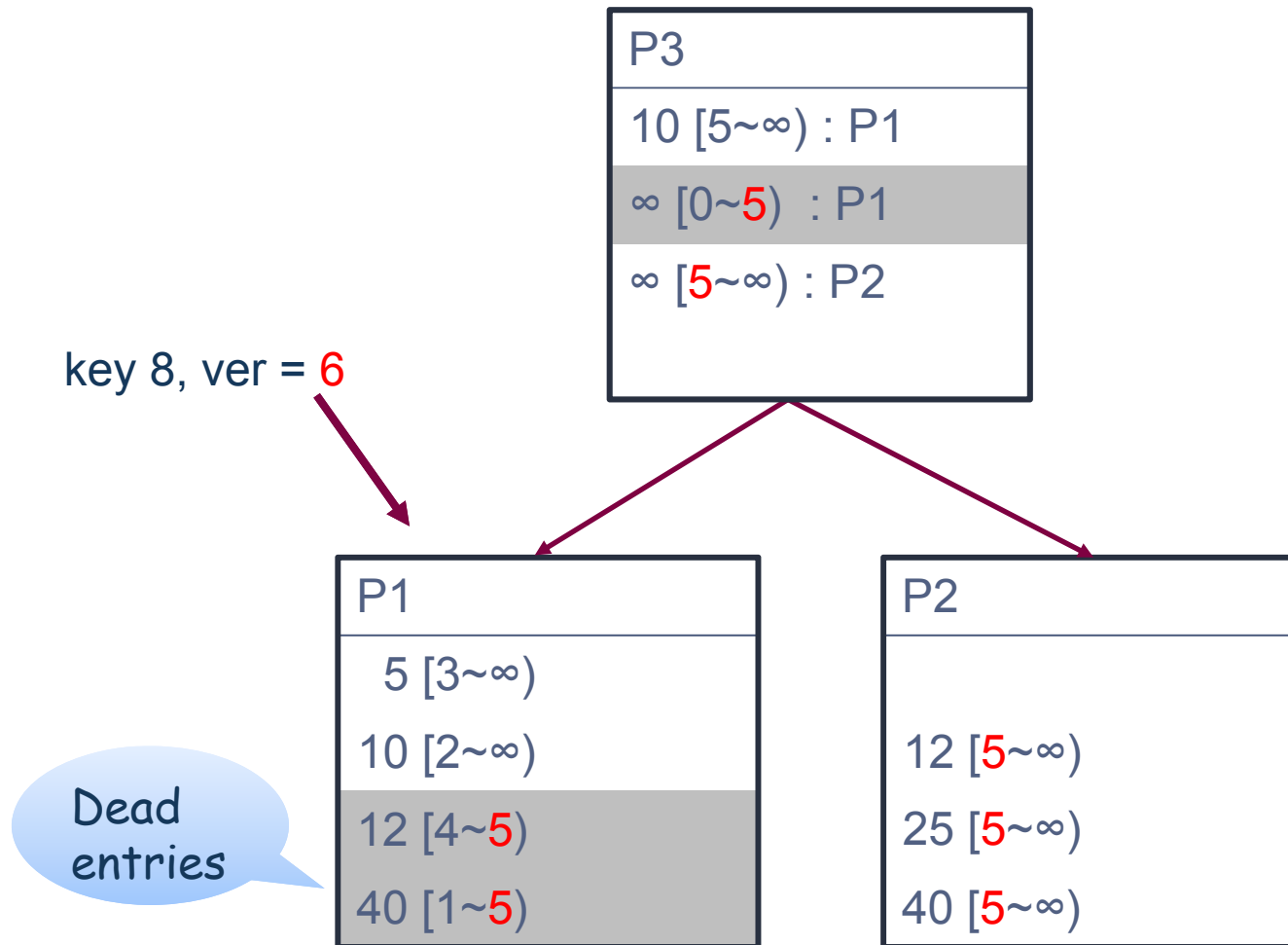
P1
5 [3~ ∞)
10 [2~ ∞)
12 [4~5)
40 [1~5)

Optimization1: Lazy Split



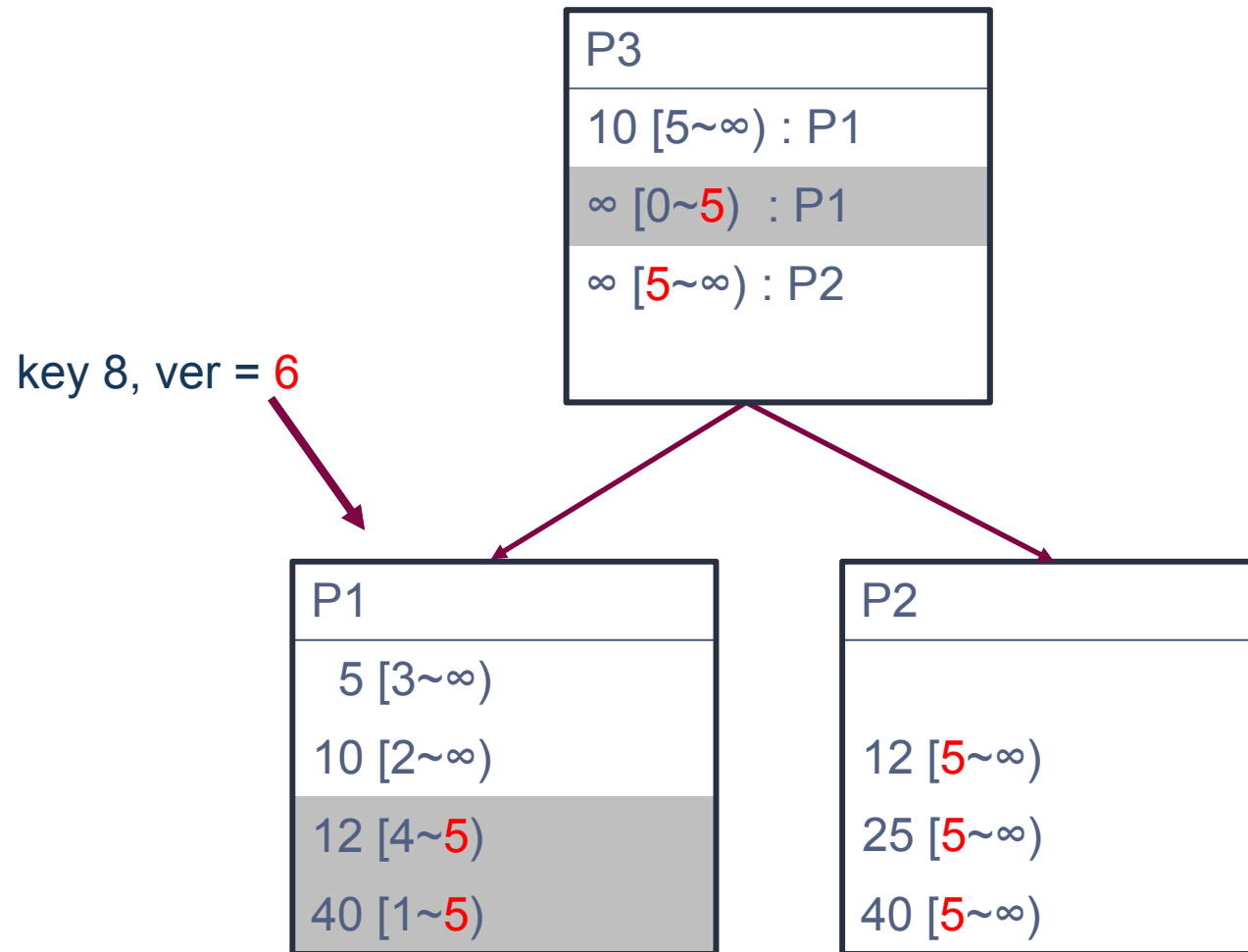
Optimization1: Lazy Split

- Lazy Node Overflow



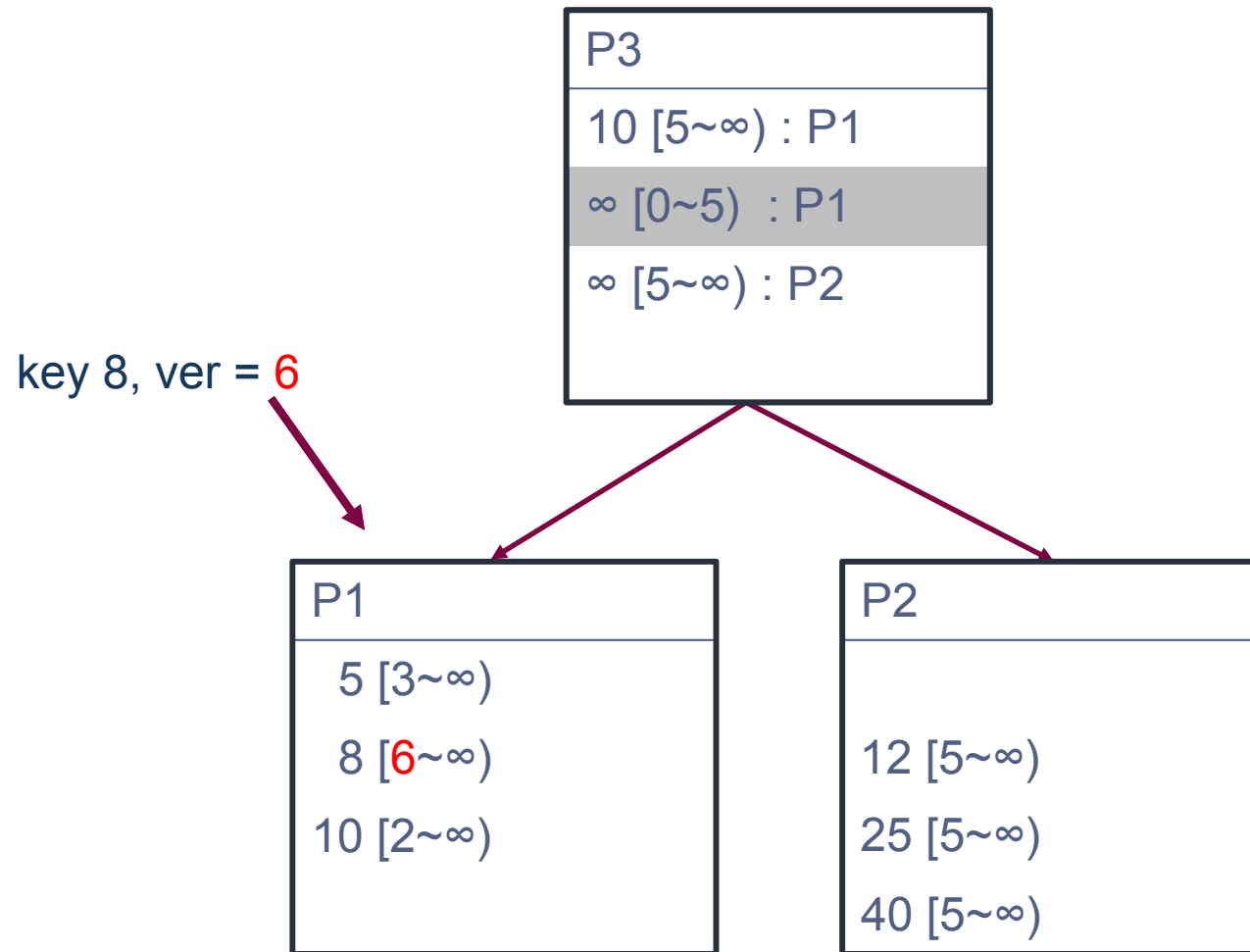
Optimization1: Lazy Split

- Lazy Node Overflow → Garbage collect dead entries



Optimization1: Lazy Split

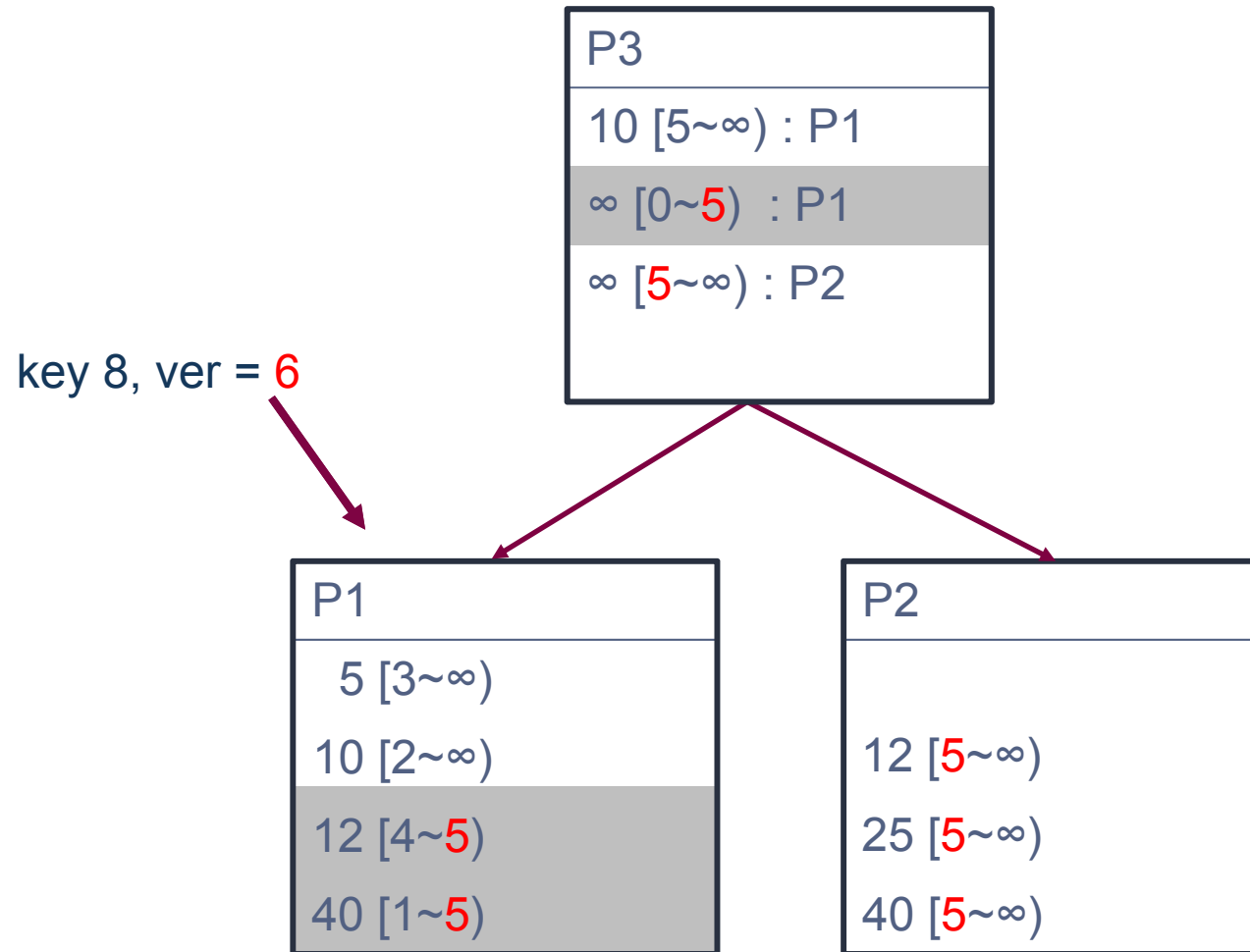
- Lazy Node Overflow → Garbage collect dead entries



But, what if the dead entries are being accessed by other transactions?

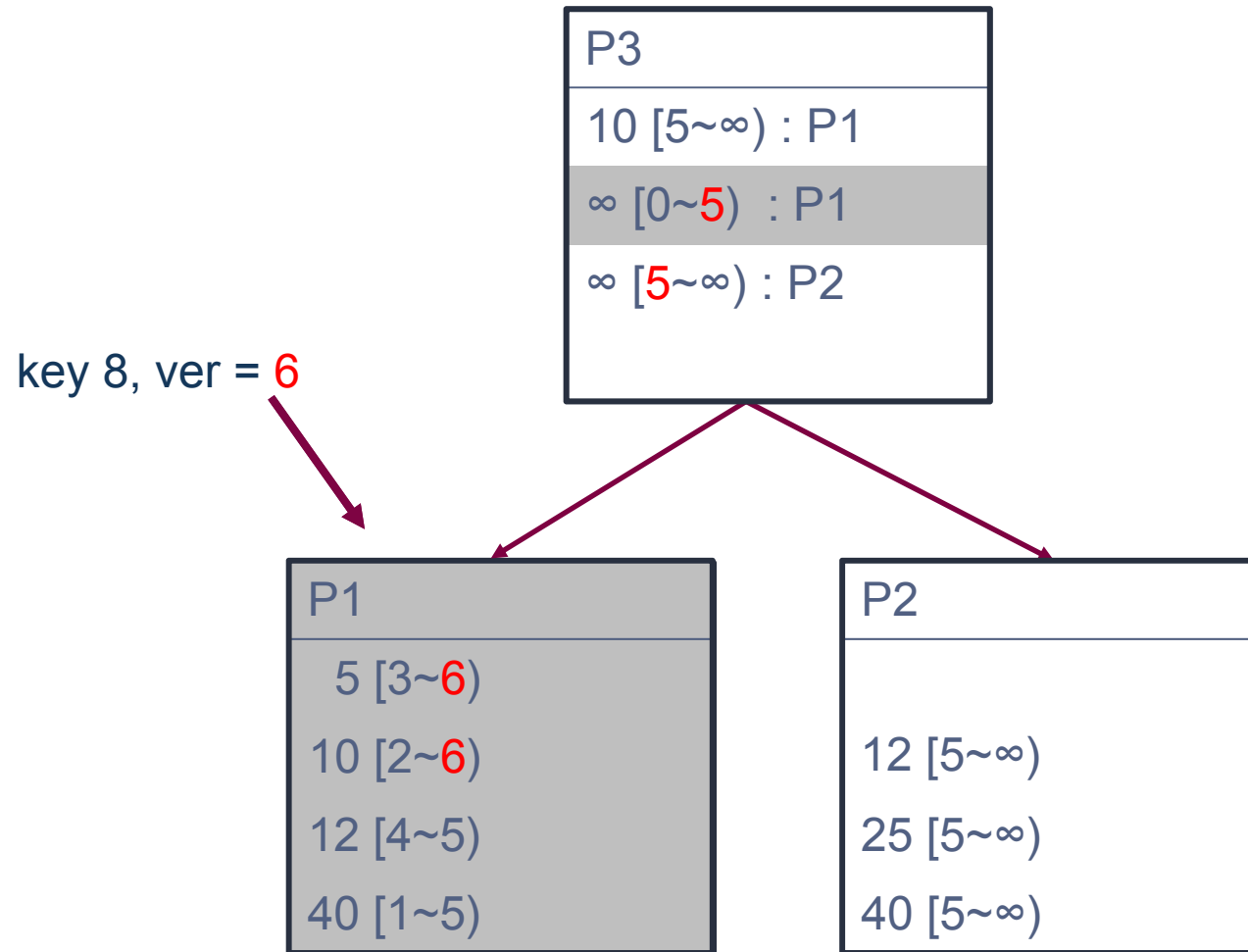
Optimization2: Reserved Buffer Space

- Option 1. Wait for read transactions to finish



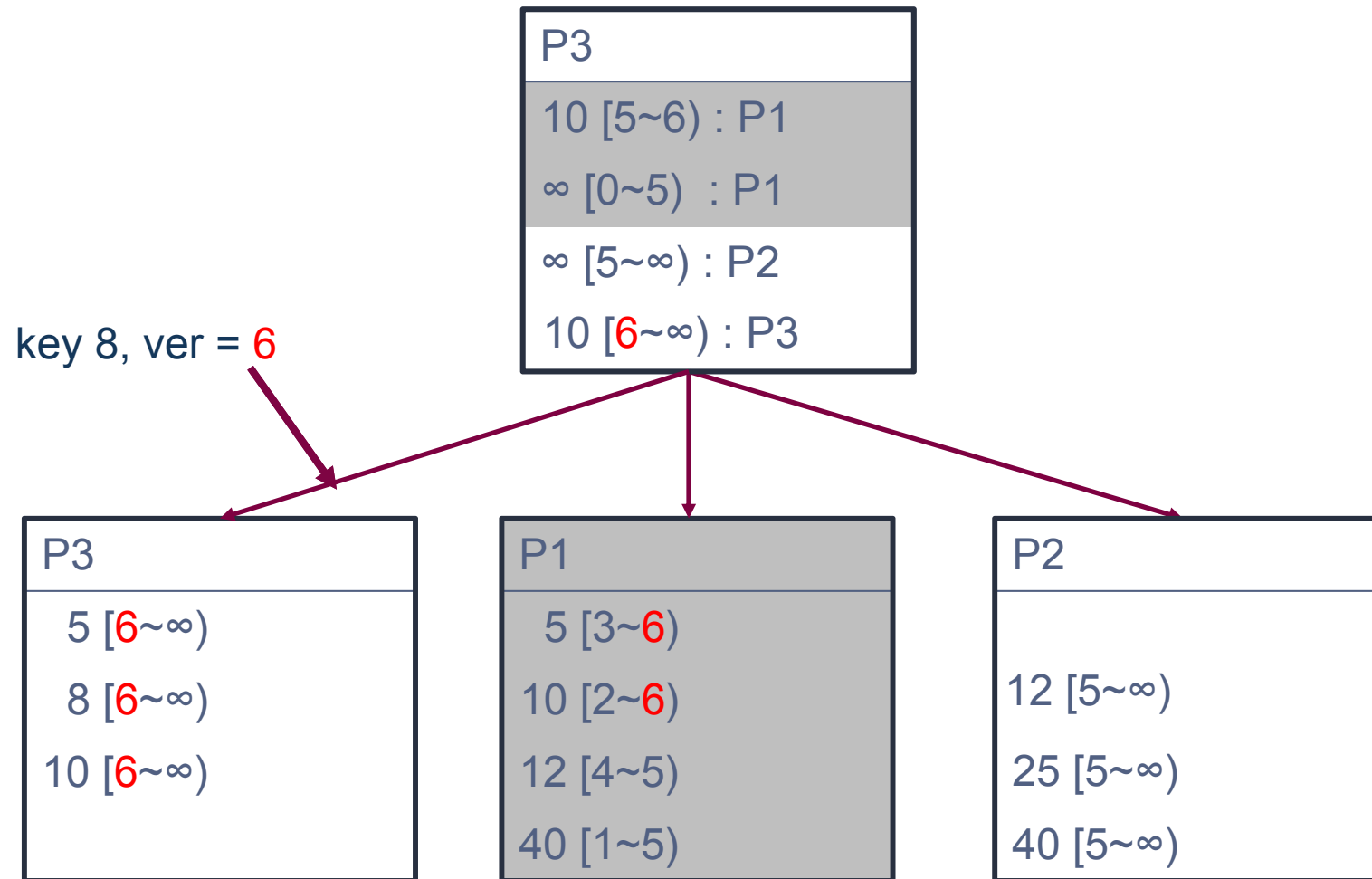
Optimization2: Reserved Buffer Space

- Option 2. Split as in legacy MVBT split



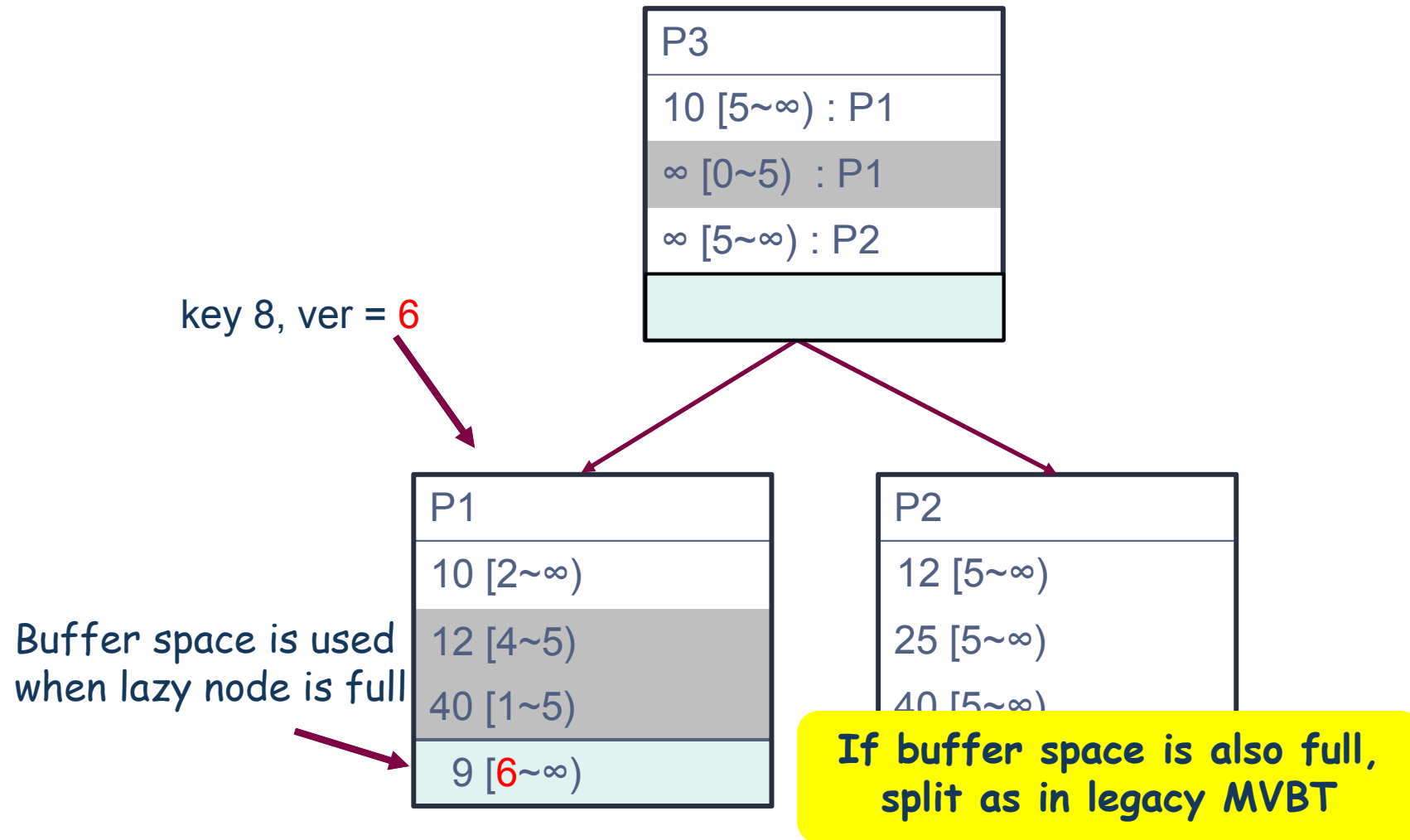
Optimization2: Reserved Buffer Space

- Option 2. Split as in legacy MVBT split



Optimization2: Reserved Buffer Space

- Option 3. Pad some buffer space in tree nodes

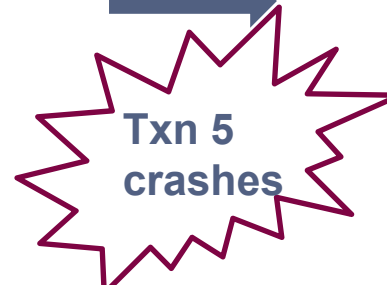


Rollback in LS-MVBT

- Similar to rollback in MVBT

P3
10 [5~∞) : P1
∞ [0~5) : P1
∞ [5~∞) : P2

Rollback



P3
∞ [0~∞) P1

Number of dirty nodes touched by rollback of LS-MVBT is also smaller than that of MVBT.

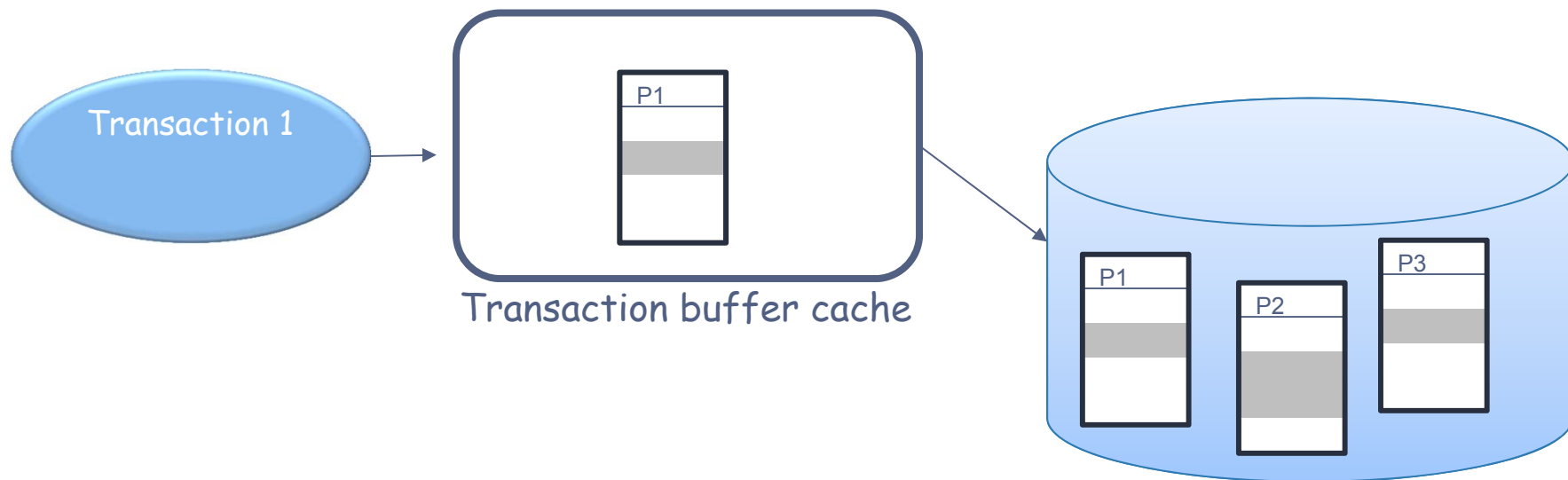
P1
5 [3~∞)
10 [2~∞)
12 [4~5)
40 [1~5)

P2
12 [5~∞)
25 [5~∞)
40 [5~∞)

P1
5 [3~∞)
10 [2~∞)
12 [4~∞)
40 [1~∞)

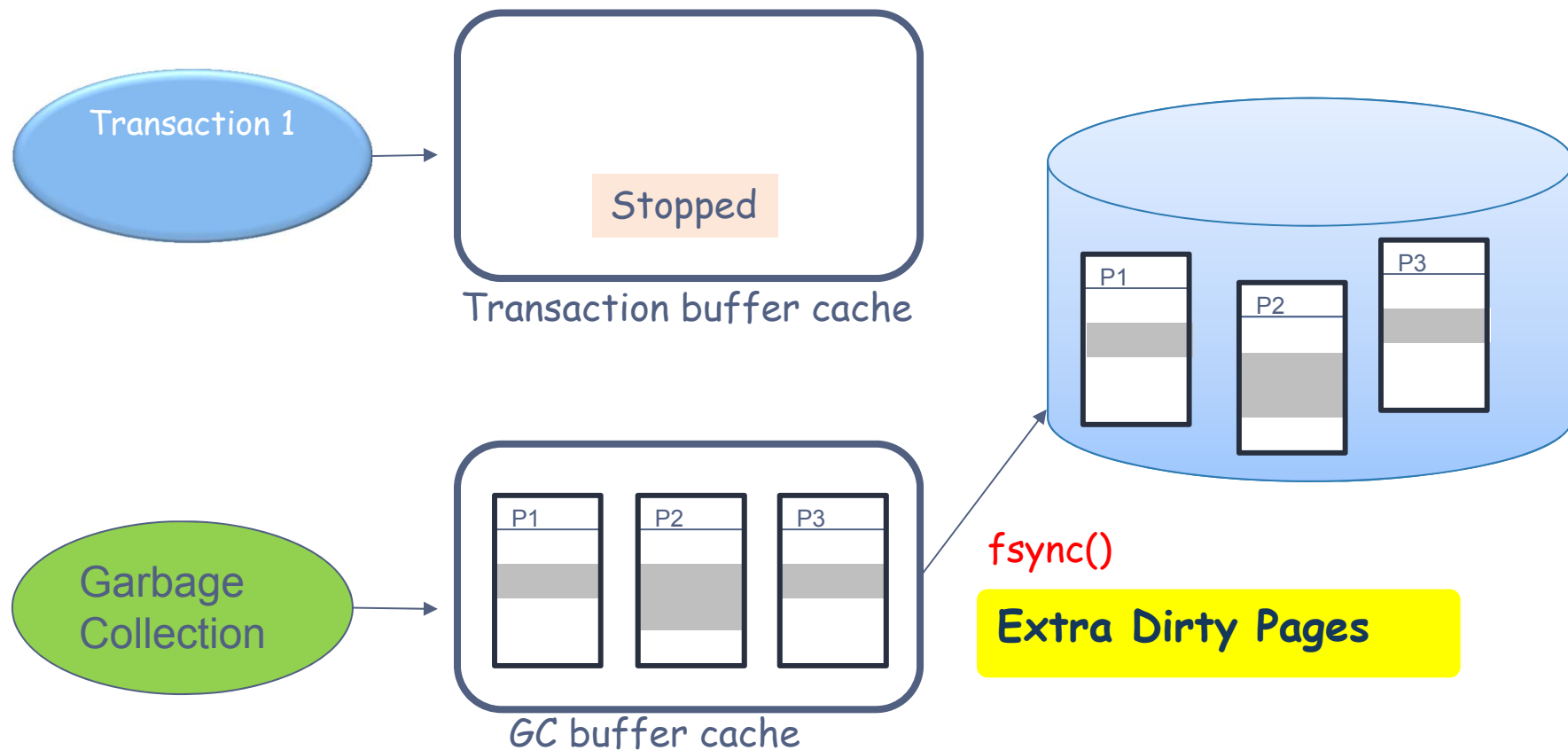
Optimization3: Lazy Garbage Collection

- Periodic Garbage Collection



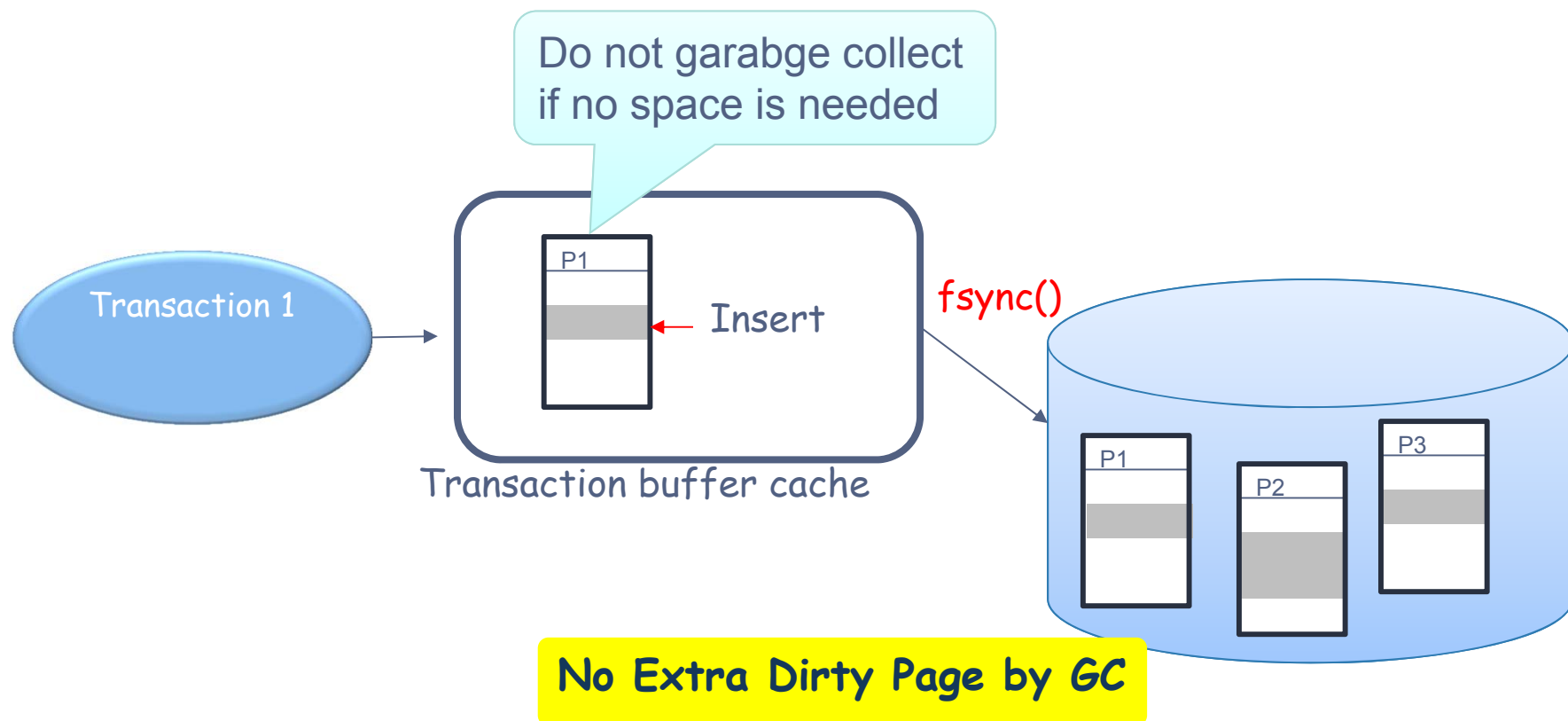
Optimization3: Lazy Garbage Collection

- Periodic Garbage Collection



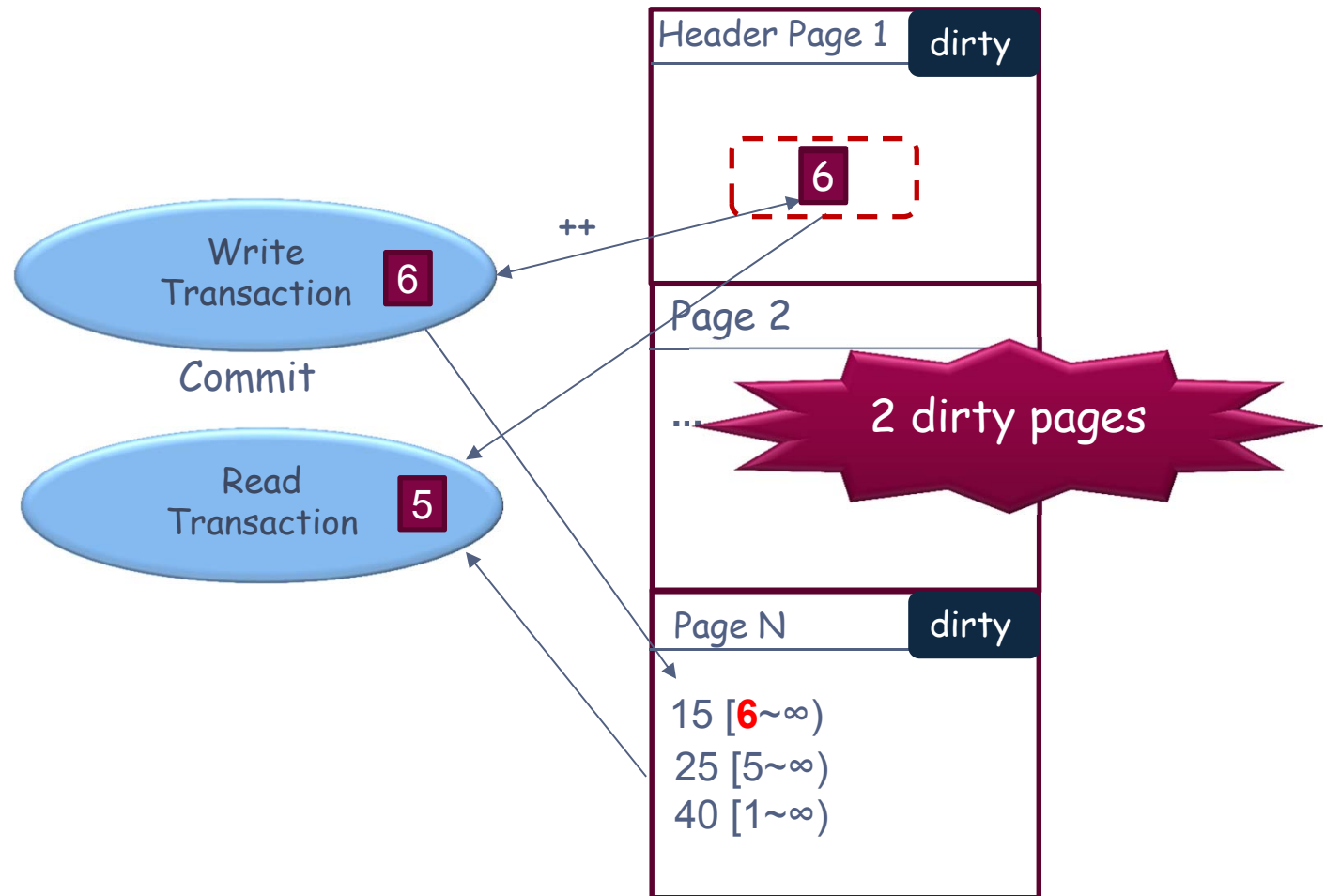
Optimization3: Lazy Garbage Collection

- Lazy Garbage Collection



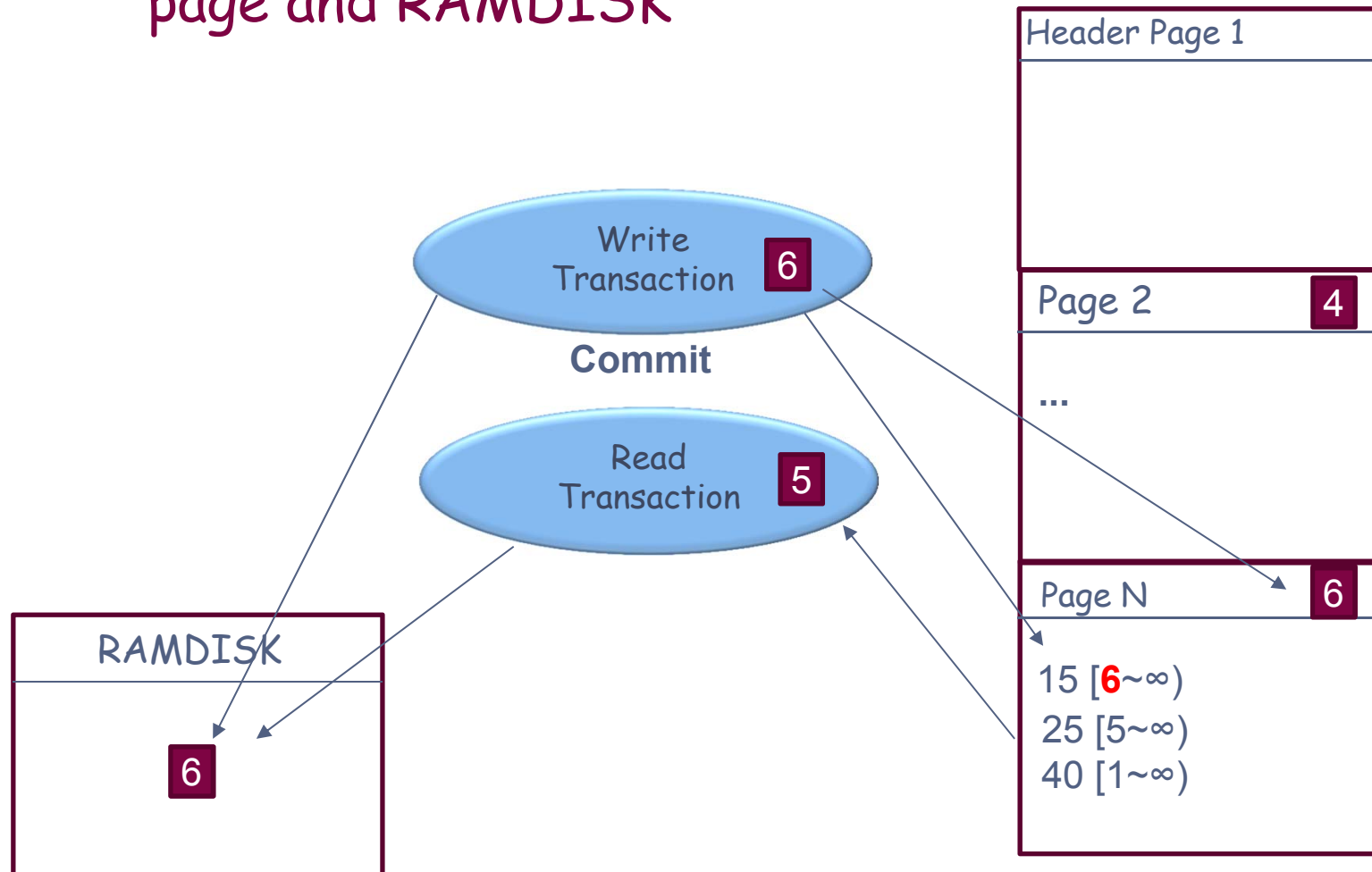
Optimization4: Metadata embedding

- Version = "File Change Counter" in DB header page



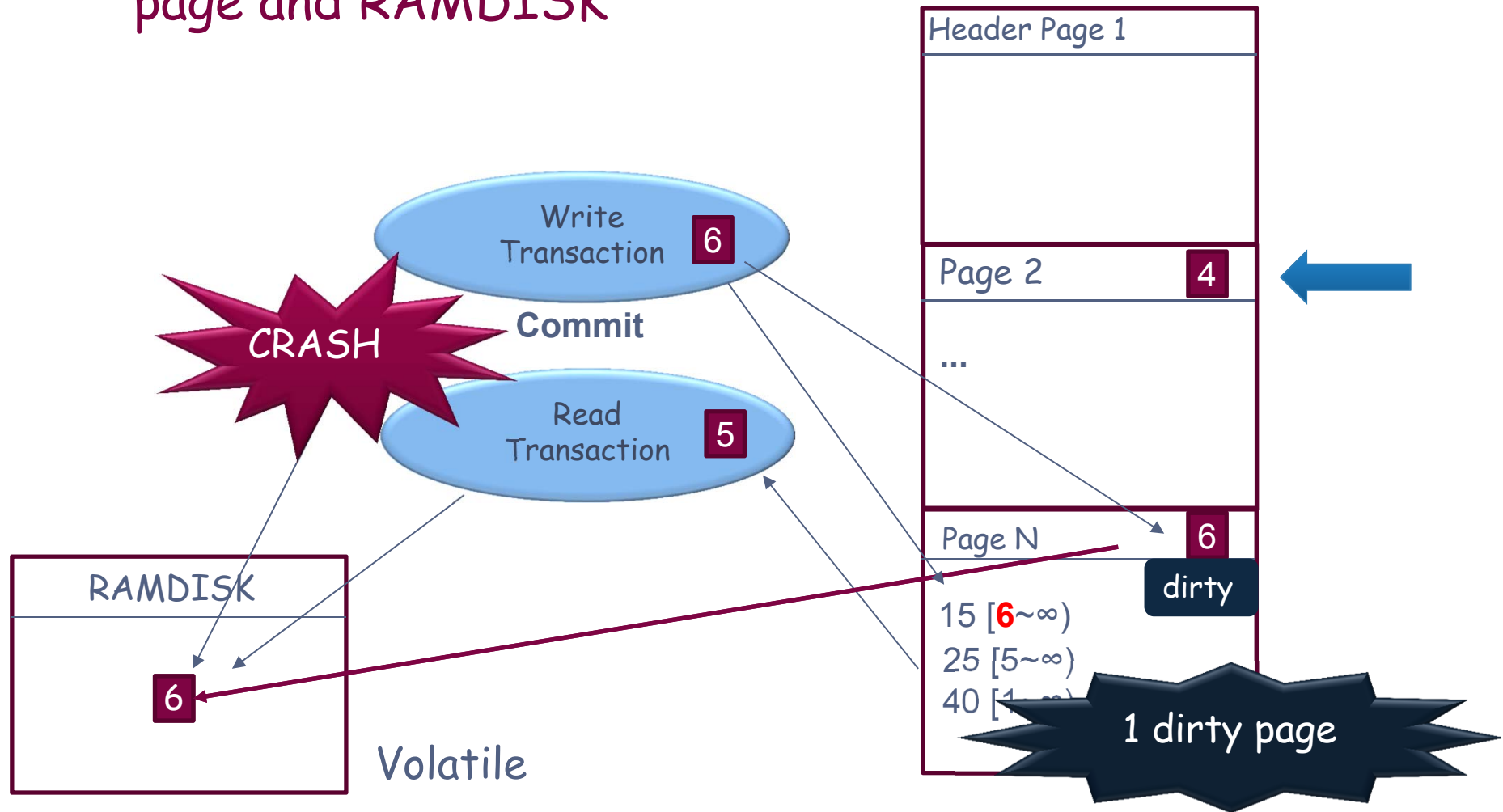
Optimization4: Metadata embedding

- Flush "File Change Counter" to the last modified page and RAMDISK



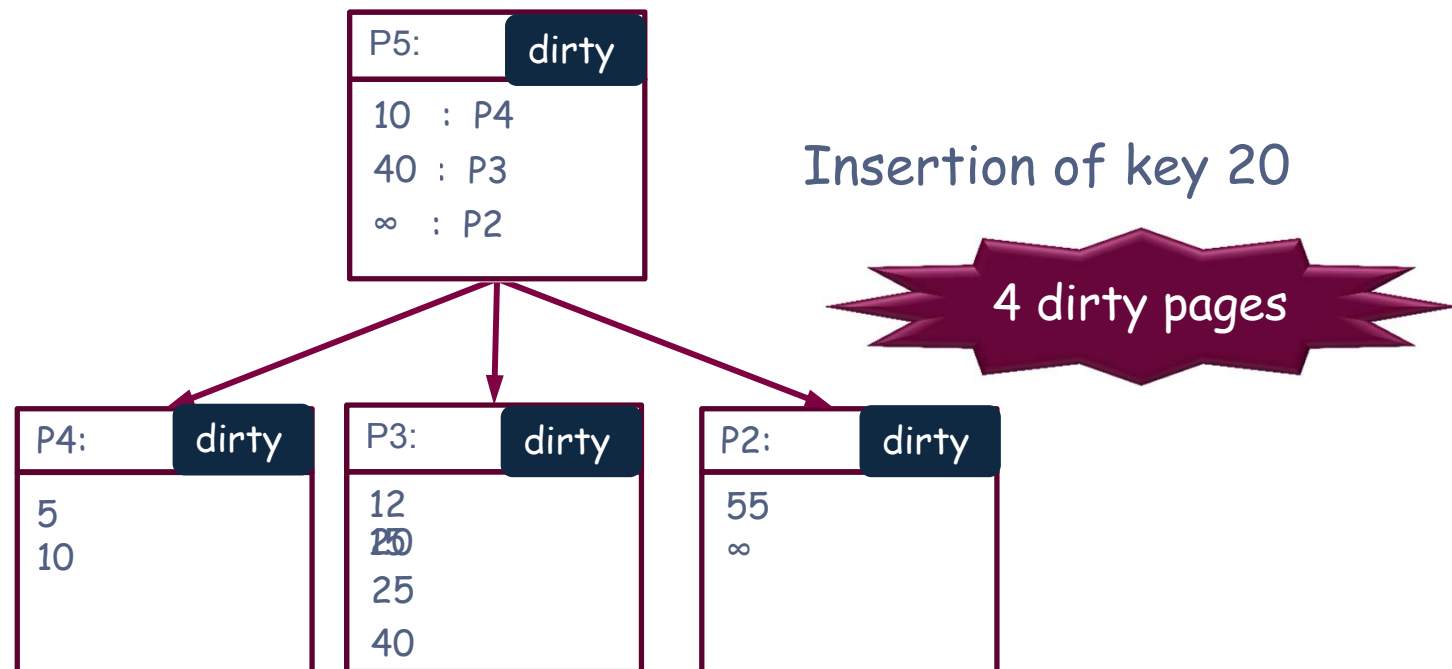
Optimization4: Metadata embedding

- Flush "File Change Counter" to the last modified page and RAMDISK



Optimization5: Disable Sibling Redistribution

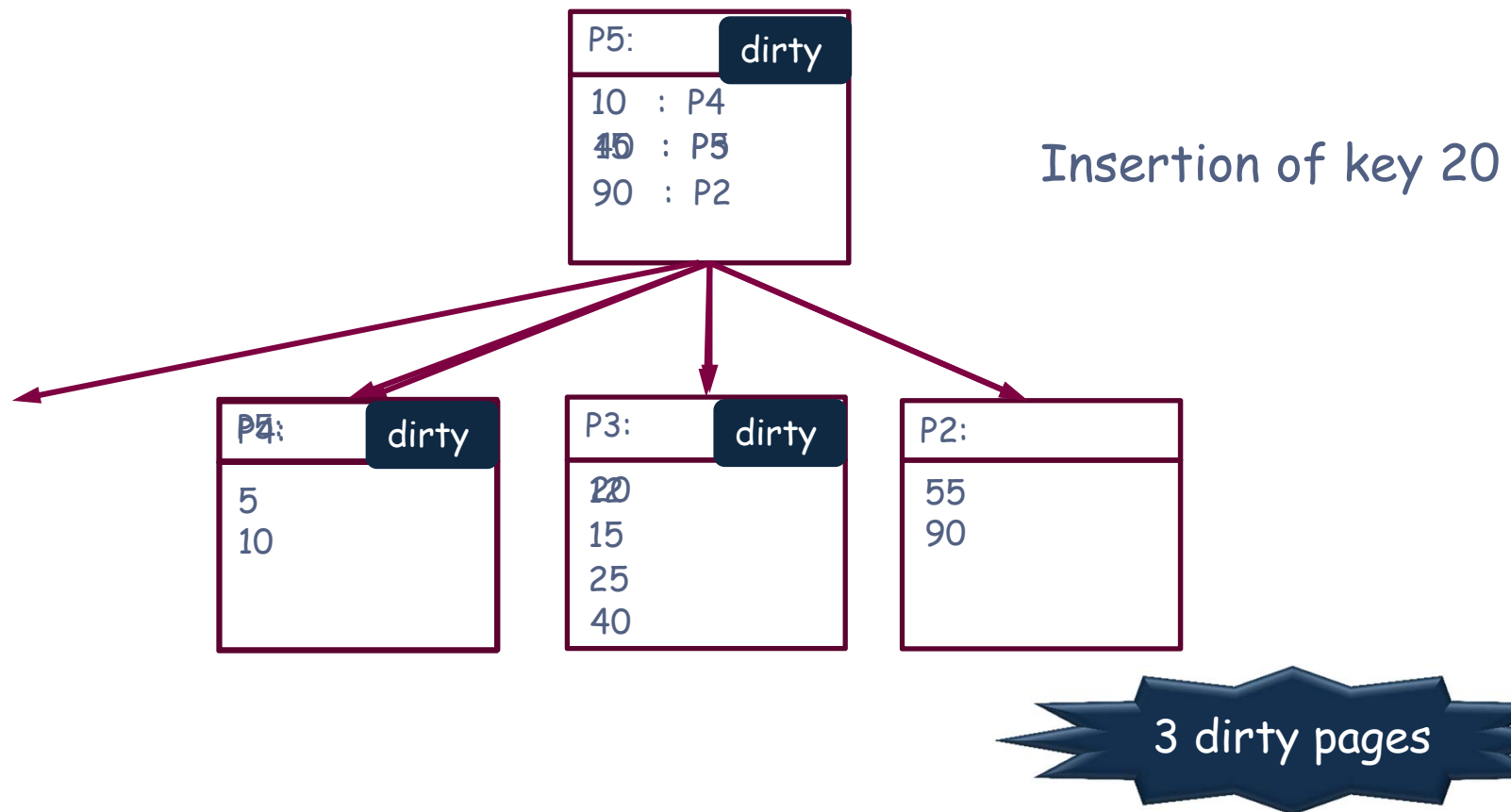
- Sibling redistribution hurts insertion performance



- Disable sibling redistribution → avoid dirtying 4 pages

Optimization5: Disable Sibling Redistribution

- Disabled sibling redistribution



Summary

■ Optimizations



Performance Evaluation

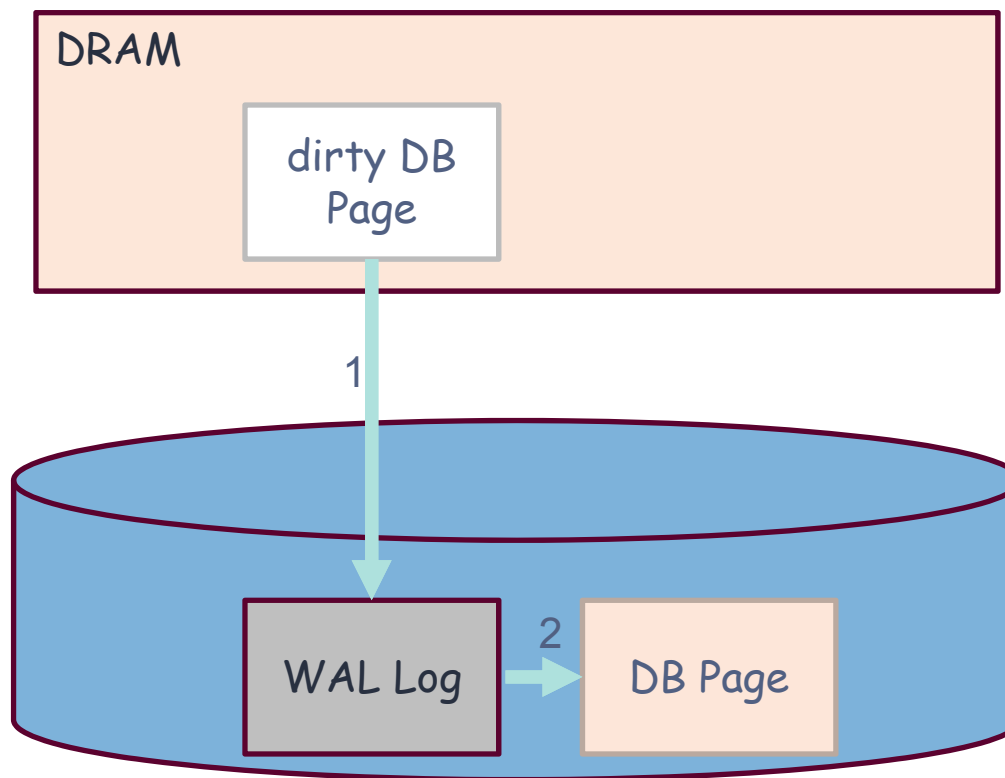
- Implemented LS-MVBT in SQLite 3.7.12
- Testbed: Samsung Galaxy-S4 (JellyBean 4.3)
 - Exynos 5 Octa 5410 Quad Core 1.6GHz
 - 2GB Memory
 - 16GB eMMC
- Performance Analysis Tools
 - Mobibench
 - MOST

Samsung
GALAXY S4



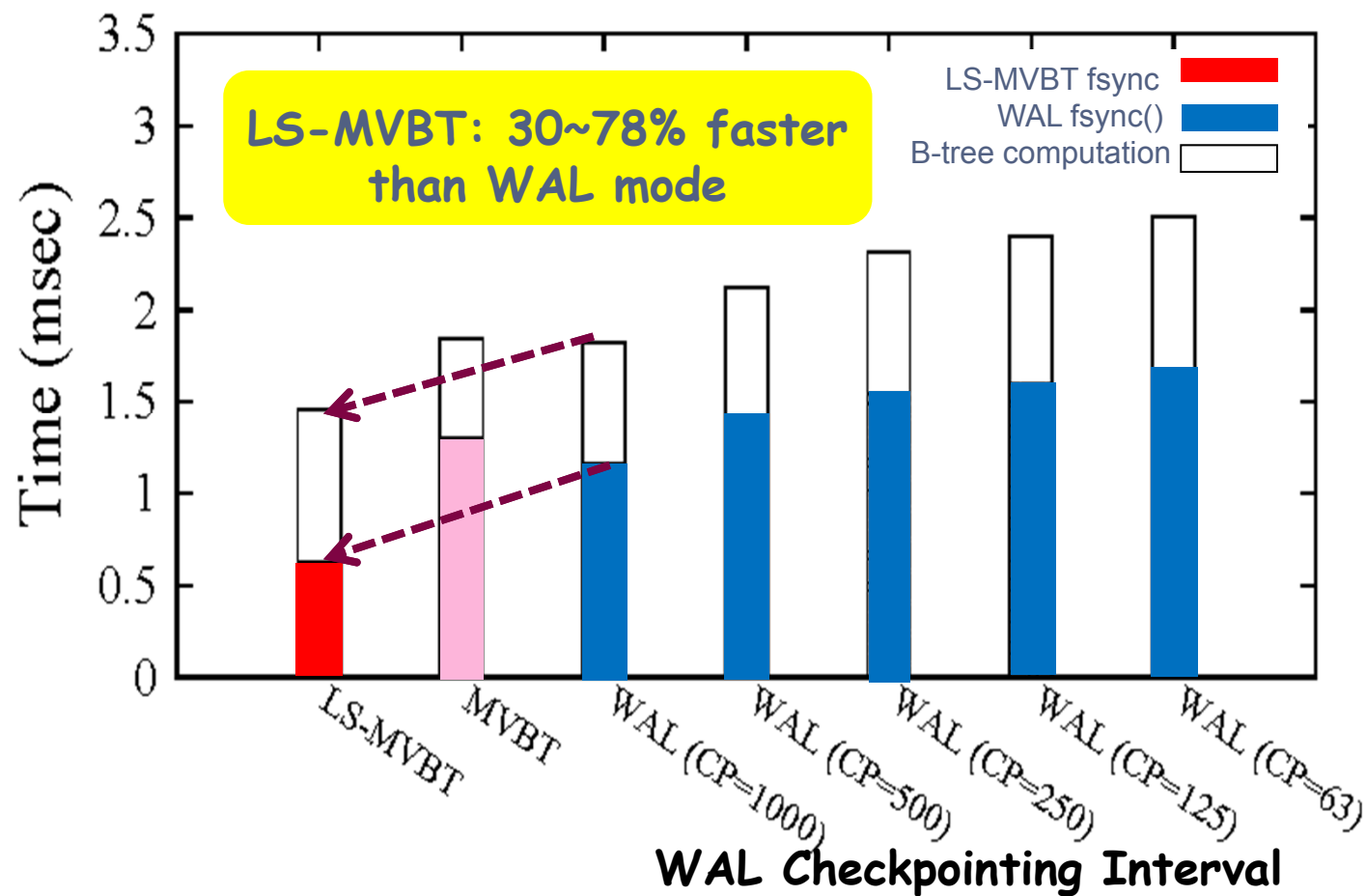
WAL (Write-Ahead-Logging) mode

- Default mode in Jelly Bean and KitKat
- Commit -> Write to a log file
- Checkpointing -> Flush logs to DB pages



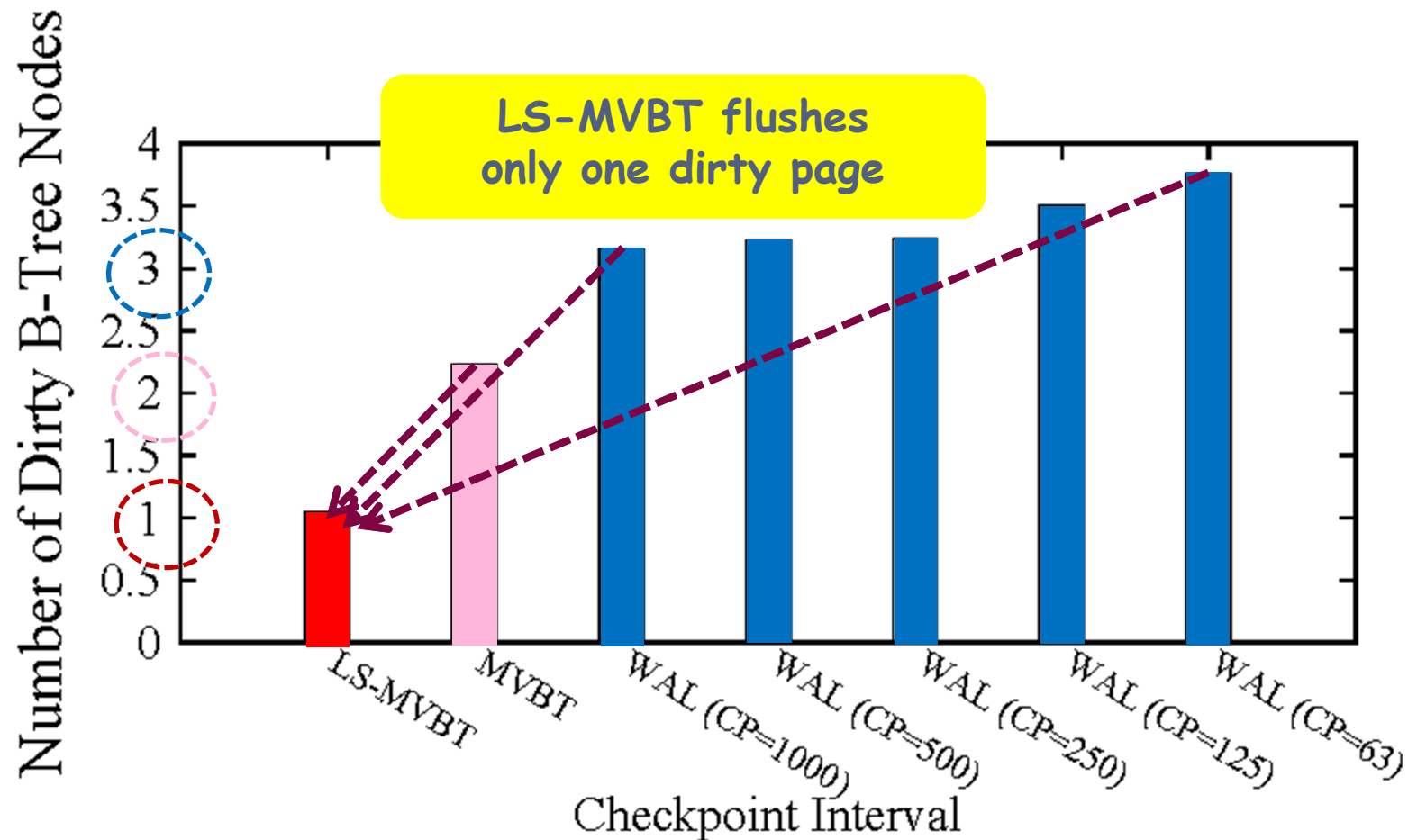
Analysis of insert Performance

SQLite Insert (Response Time)



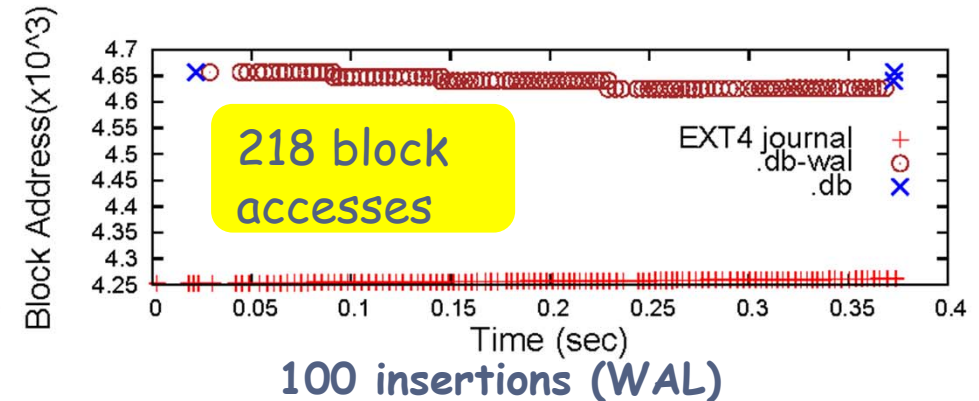
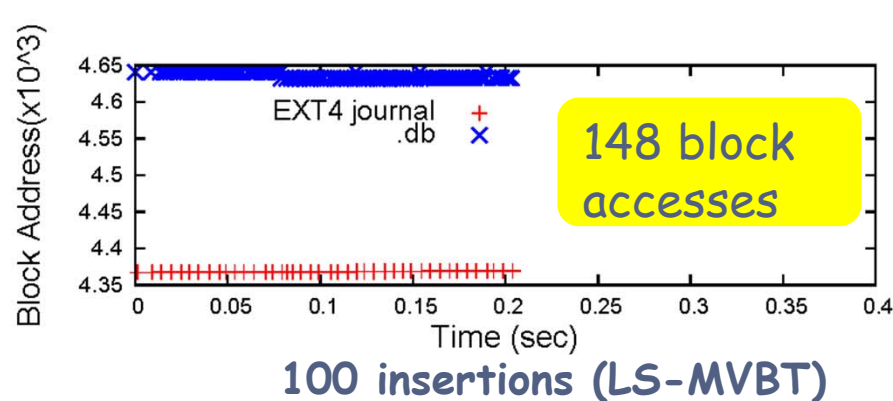
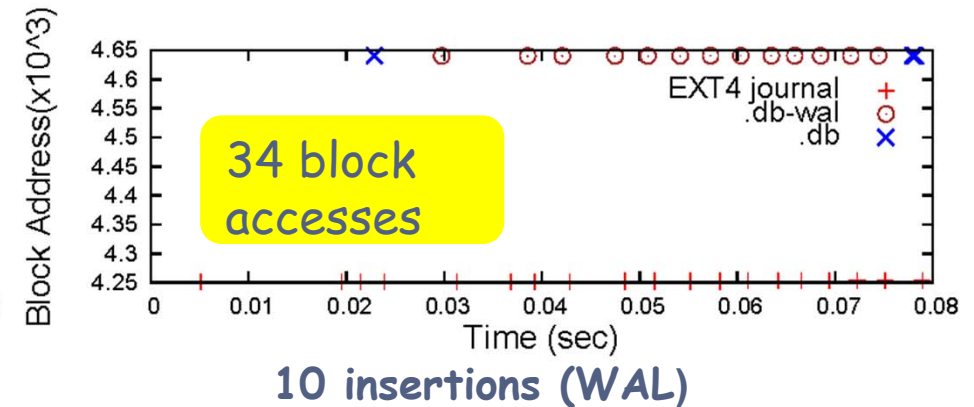
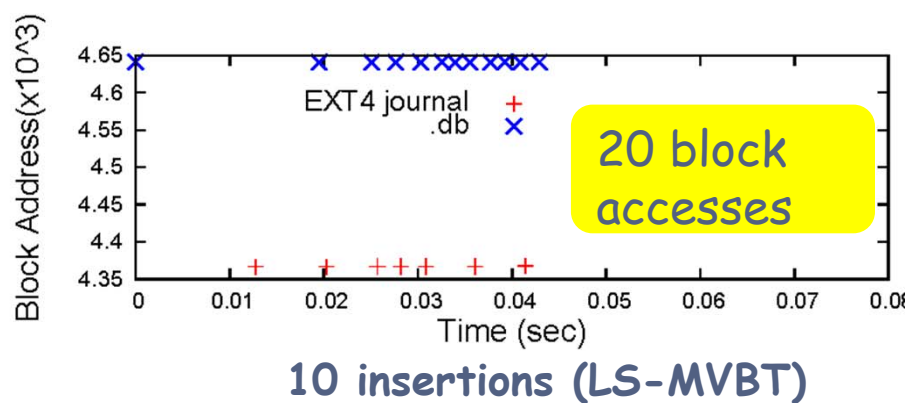
Analysis of insert Performance

SQLite Insert (# of Dirty Pages per Insert)



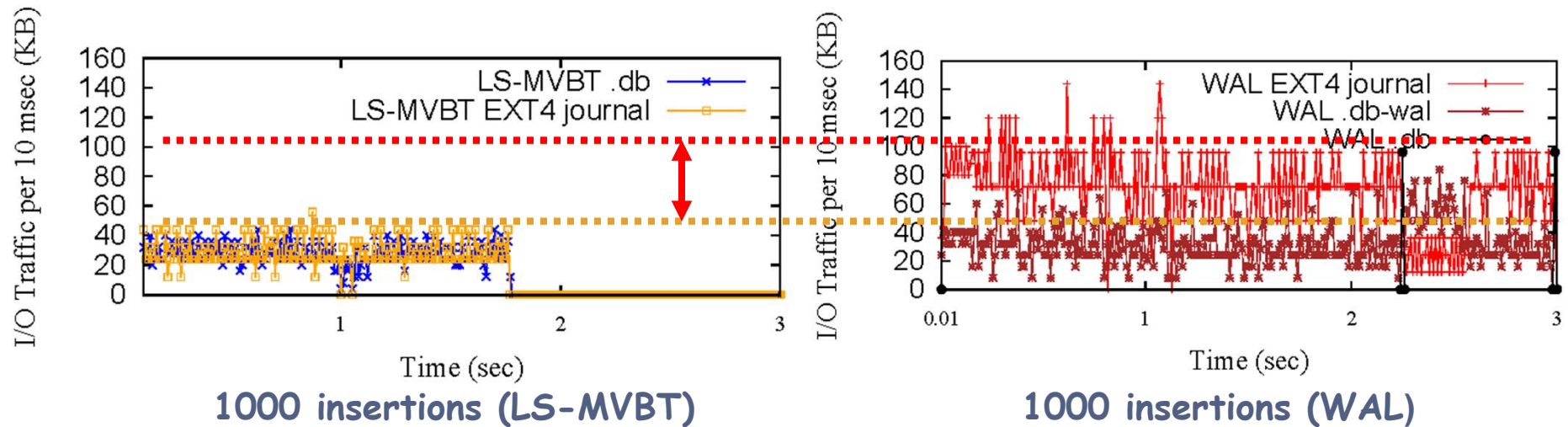
Analysis of Block I/O Behavior

SQLite Insert (# of Accesses to Block Device)



I/O Traffic at Block Device Level

SQLite Insert (I/O Traffic per 10 msec)



LS-MVBT saves 67% I/O traffic:
9.9 MB (LS-MVBT) vs 31 MB (WAL)

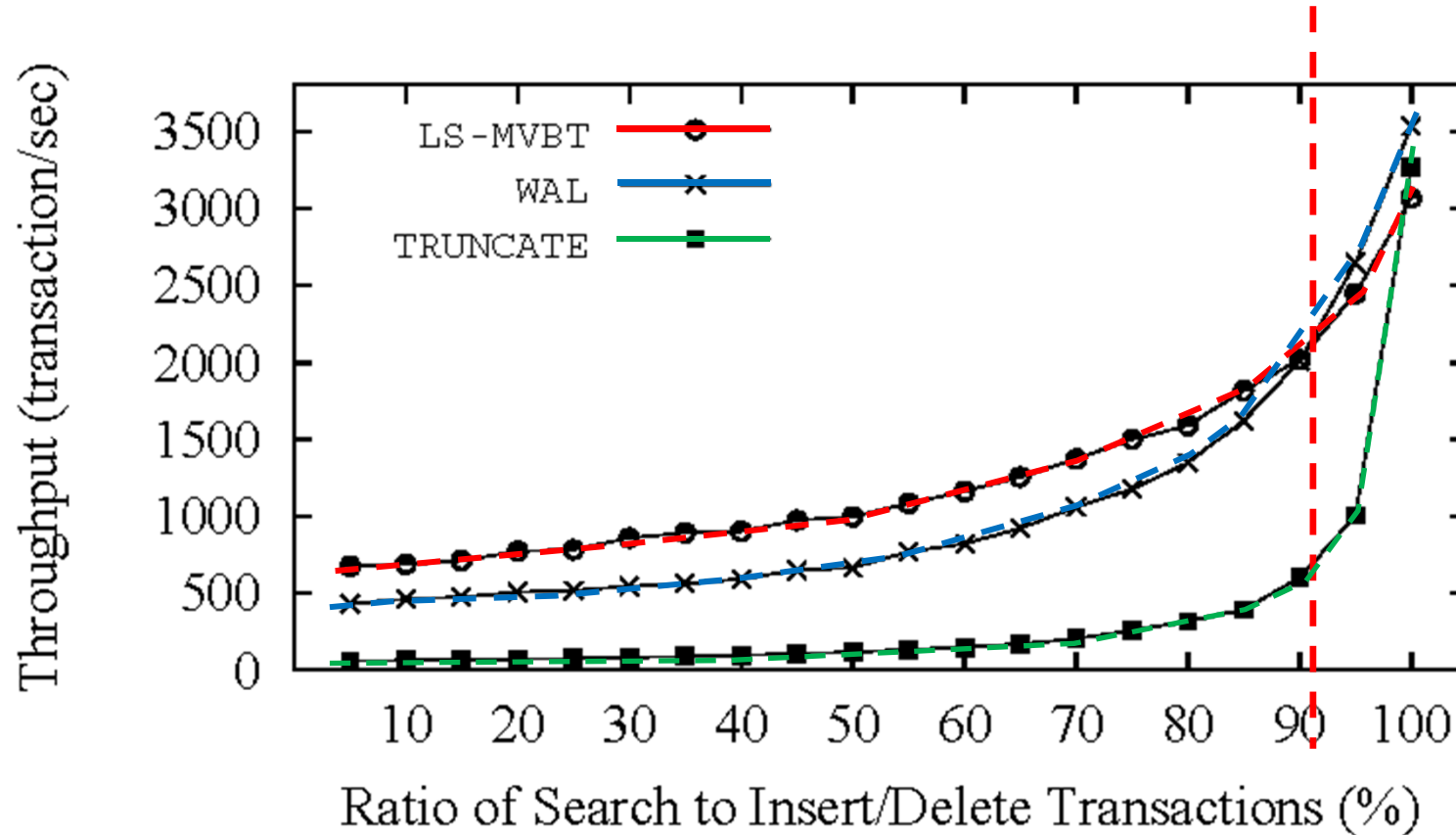
x3 longer life time of NAND flash



- How about Search performance?
 - LS-MVBT is optimized for write performance.

Search Performance

Transaction Throughput with varying Search/Insert ratio



LS-MVBT wins unless more than 93% of transactions are search

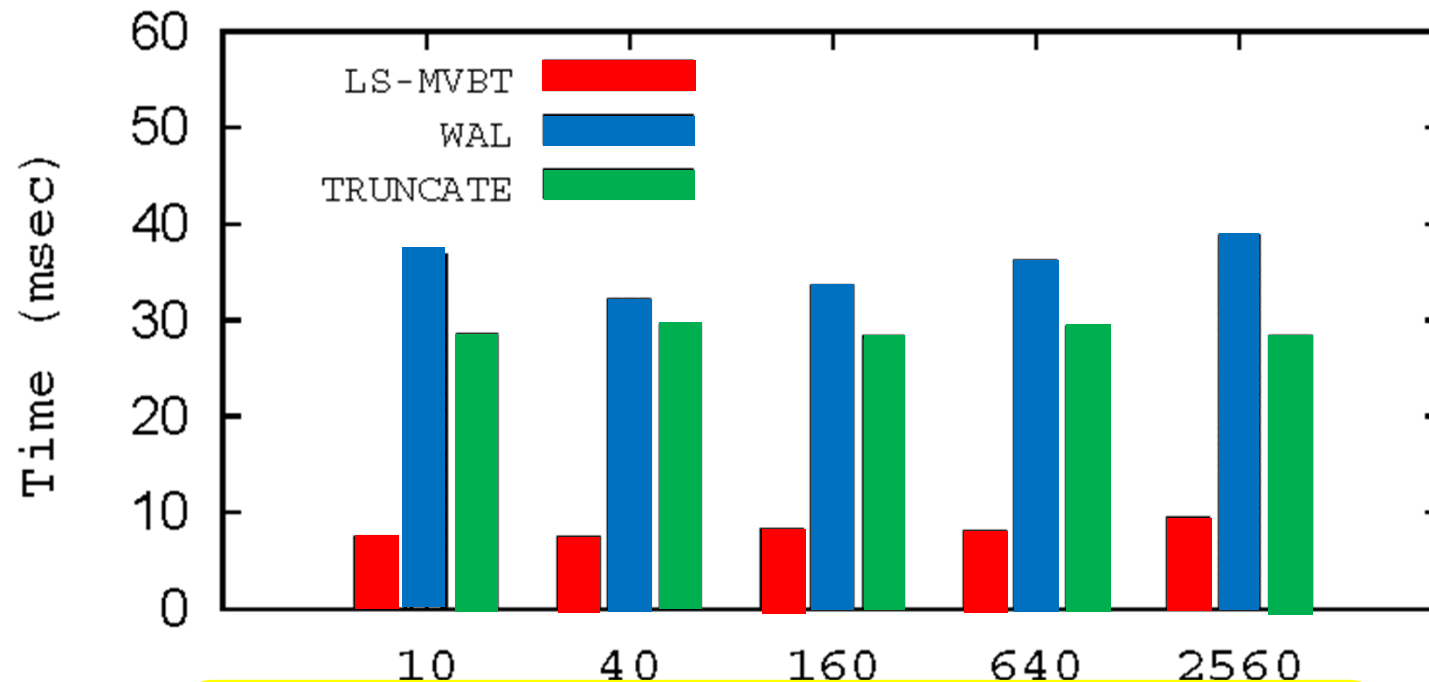


- How about recovery latency?

- WAL mode exhibits longer recovery latency due to log replay.
- How about LS-MVBT?

Recovery

Recovery Time with varying # of insert operations to rollback



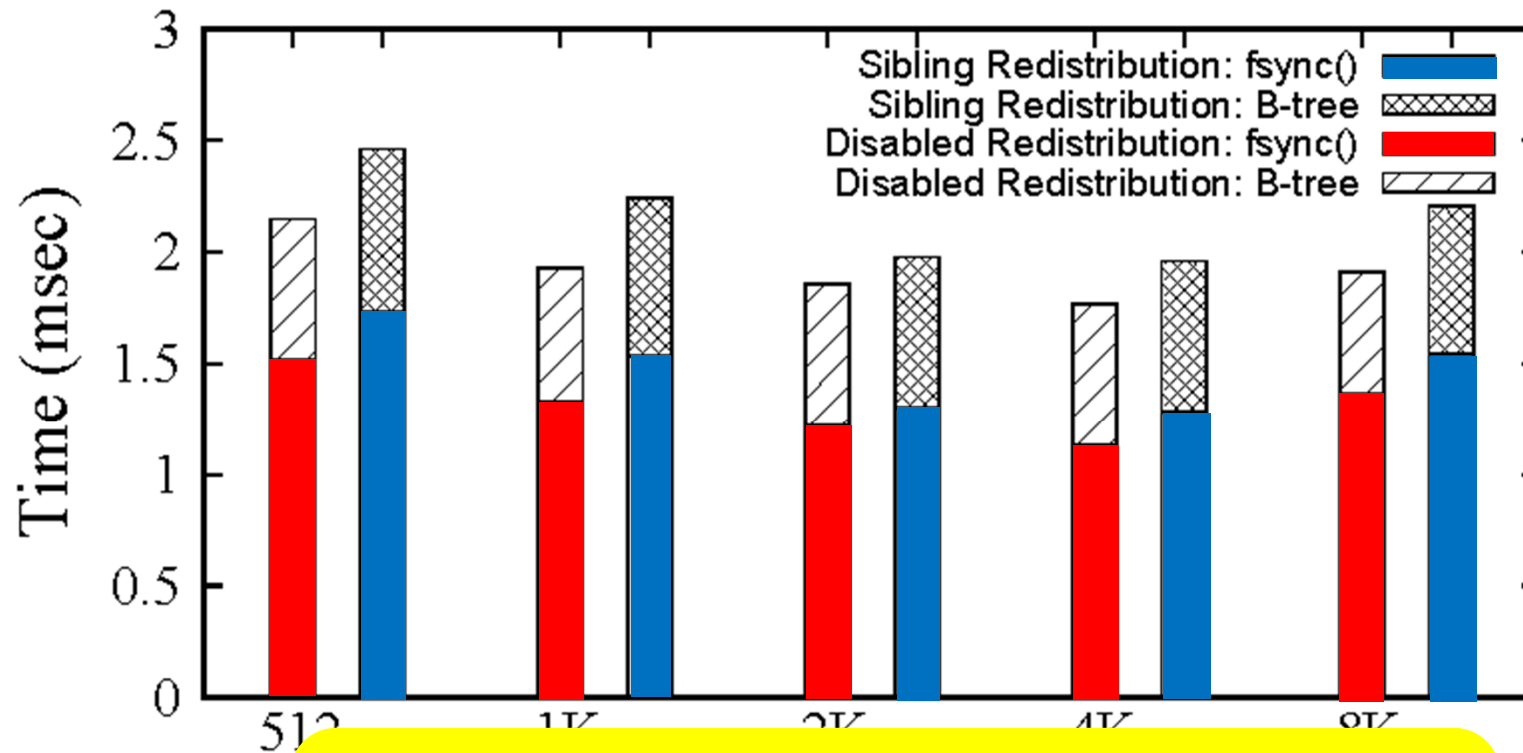
LS-MVBT is **x5~x6** faster than WAL.



- How much performance improvement for each optimization?

Effect of Disabling Sibling Redistribution

SQLite Insert: (Response Time and # of Dirty Pages)



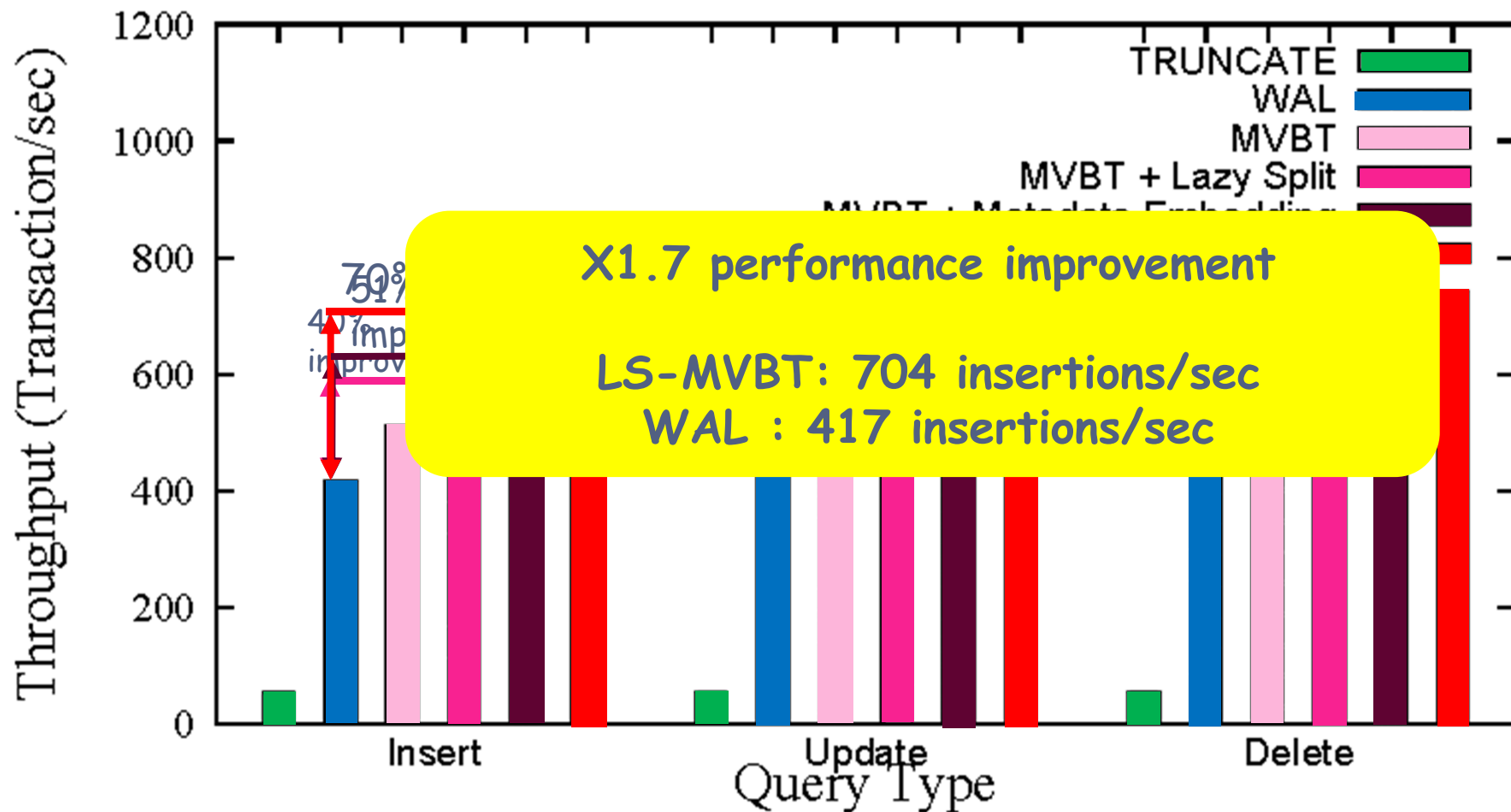
Disabling sibling redistribution



20% faster insertion

Performance Quantification

Throughput of SQLite: Combining All Optimizations



Conclusion

- LS-MVBT resolves the Journaling of Journal anomaly via
 - Reducing the number of fsync()'s with **Version based B-tree**.
 - Reducing the overhead of single fsync() via **lazy split**, **disabling sibling redistribution**, **metadata embedding**, **reserved buffer space**, **lazy garbage collection**.
- LS-MVBT achieves **70% performance** gain against WAL mode (417 ins/sec → 704 ins/sec) **solely via software optimization**.

Future/Ongoing Works

1. Improve the performance of WAL mode
 - 3 dirty pages per insertion → not really necessary
 - Easier work than replacing the entire B-tree
 - Compatible with current SQLite database files
 - We expect the performance benefit would be similar to LS-MVBT
2. SQLite for Non-Volatile Memory
 - Logging/Journaling in NVRAM
3. SQLite for multi-cores

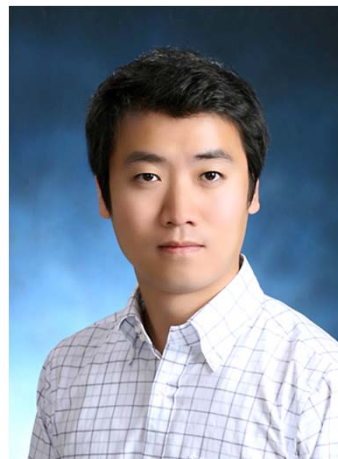
Credit



Woo-Hee Kim



Beomseok Nam



Dongil Park



Youjip Won

Thank you...

